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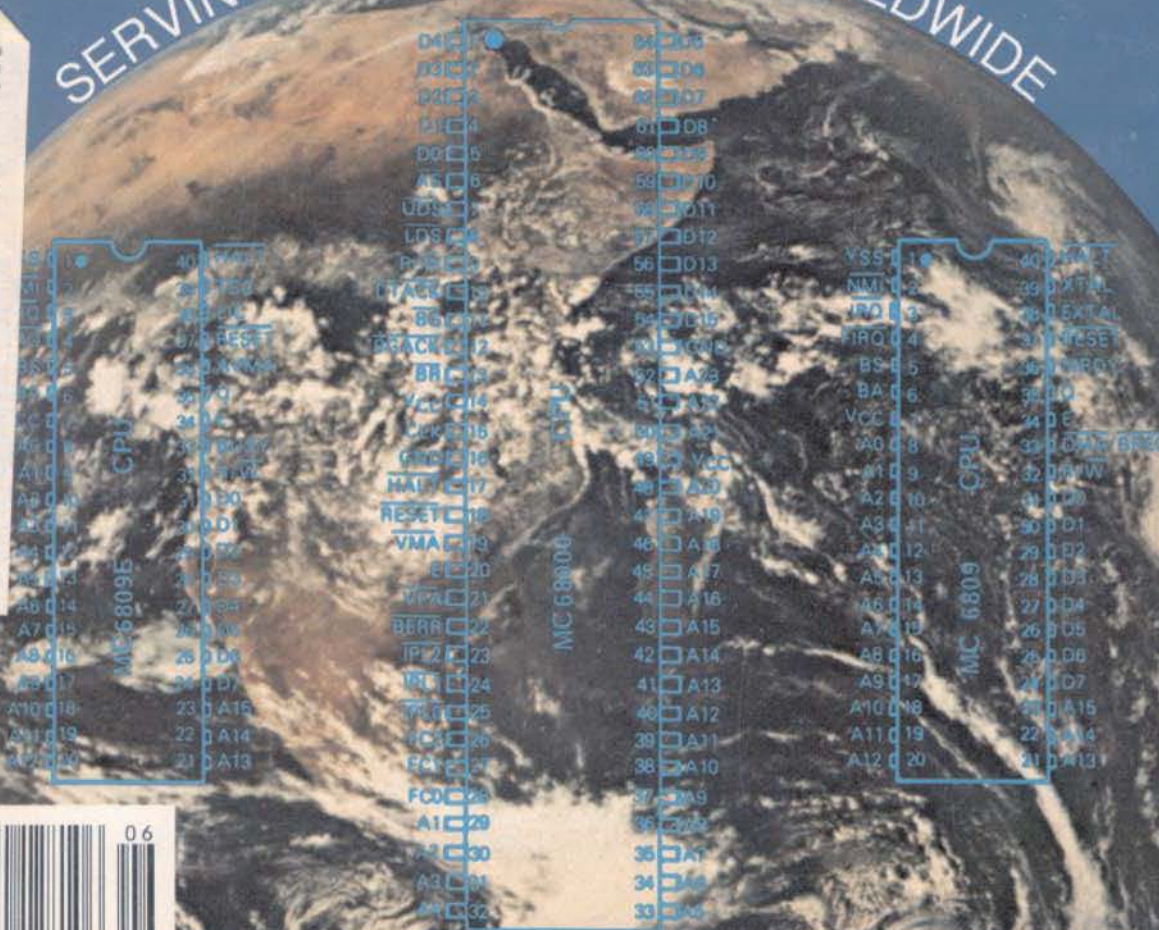
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INTRODUCTION

This chapter presents still another version of a routine to scan command lines and two utility programs which are heavily dependent upon this scanner. They were all written by John Weald and placed into the public domain for the rest of us to use. The scanner is a version of getopt, a UNIX function, and the utility programs are versions of cut and paste, useful UNIX commands for text processing.

GENERAL

Command line scanners are not difficult to write; however, many have been written without thought as to how they are to be used. Some seem to have been designed to be difficult or obscure in their use, for no apparent reason.

Ideally, the order of options with respect to other arguments, such as file names, should logically be irrelevant in most programs. Thus, the following sequences would then generate equivalent results:

```
prog -a -b -c file1 file2 file3
prog file1 file2 file3 -a -b -c
prog -a file1 -b file2 -c file3
```

Such a capability is appreciated by many who use computers heavily and do not wish to consult manuals concerning the order of options. If order dependencies are present in a given program, they should be documented in a usage message (all programs have good usage messages, don't they?) or in a manual.

Many commonly-used programs on popular operating systems have such dependencies, and it is left to the programmer to remember the requirements and peculiarities of each program. The really treacherous programs accept the options without error but do not act on them as expected.

Microsoft C version 5.0 command-line arguments are scanned in a left to right manner. Thus, in the following example,

```
cl prog1.c -Dmsc -Am
```

the -Dmsc and -Am options would not be processed until after prog1.c is processed, potentially leading to a program being silently compiled with incorrect options.

The getopt scanner below provides the same capability as one by a similar name defined under UNIX. Unfortunately, it is not an ideal scanner in that it requires all options to be coded before other arguments, unless the programmer makes the extra effort to call it at appropriate times. Remember to provide an indicative usage message in such cases.

GETOPT

Getopt is declared with the following formal parameters:

```
int getopt (argc, argv, optstring)
int argc;
char **argv, *optstring;
```

It has the following external references:

```
extern char *optarg;
extern int optind, opterr;
```

Getopt returns the next option letter in argv (starting from argv[1]) which matches a letter in optstring, which is a string of recognized option letters; if a letter is followed by a colon, the option is expected to have an argument that may or may not be separated from it. Optarg is set to point to the start of the option argument on return.

Getopt places in optind the argv index of the next argument to be processed. Optind is initialized to 1 before the first call to the getopt function.

Options can be any ASCII characters except colon, question mark, or null. When all options have been processed, getopt returns EOF. The special option — may be used to delimit the end of the options; EOF will be returned, and — will be ignored.

Getopt normally outputs an error message to stderr and returns a question mark when it encounters an option letter not included in optstring. It is impossible to distinguish between a ? used as a legal option, and the character that getopt returns when it encounters an invalid option character in the input. This error message may be suppressed by setting opterr to zero.

```c
/*
 *    This is a clone of the UNIX getopt() function.
 *
 * John Weald
 */

#include <stdio.h>

/*
 * Index into error array.
 */
#define BAD_OPT 0 /* option letter not in optstr */
#define MIS_ARG 1 /* option must have an argument */

char *optarg;
char *p;
char errors[2][30];
int opterr = 1; /* If true print error message */
int optind = 1; /* argv[0] is program name */

/*
 * The basic data structures are optind
 * and the pointer p.  optind keeps track
 * of the next index into argv to parse
 * arguments.  p is used to walk along the
 * argv items looking for option letters or
 * arguments, when it is NULL the next argv
 * must be used.  p is always left pointing
 * to the previous option or NULL.
 * Consider the three equivalent argv's:
 *           1        2         3
 *          -a       -b       eric
 *             -ab      eric
 *             -aberic
 */

getopt(argc, argv, optstr)
int argc;
char *argv[];
char *optstr; /* The list of valid options */
{
    extern err();    /* forward reference */

    /*
     * parsed all the options in this argv[]?
     */
    if (!p || !*++p)
    {
        if (optind == argc)
            return(EOF);
        p = argv[optind];

        /* a '-' by itself is not an option */
        if (*p++ != '-' || !*p)
            return(EOF);

        /* '--' marks end of the option list */
        if (*p == '-')
        {
            optind++;
            return(EOF);
        }
    }

    optind++;
    /*
     * Look for a valid option
     */
    while (*p != *optstr)
    {
```

```c
        if (!*optstr)
        {
            /* Reached end of optstr */
            err(argv[0], BAD_OPT, *p);
            return((int)'?');
        }
        if (*++optstr == ':')
            optstr++;
    }

    /* If needs no argument we are done. */
    if (*(optstr + 1) != ':')
        return((int)*optstr);

    /*
     * If there are more characters,
     * they must be the argument.
     */
    if (*++p)
    {
        optarg = p;
        p = NULL;
        return((int)*optstr);
    }

    /*
     * It needs an argument, but have none
     */
    if (optind == argc)
    {
        err(argv[0], MIS_ARG, *optstr);
        p = NULL;
        return((int)'?');
    }

    /*
     * Must be in next argv.
     */
    optarg = argv[optind++];
    p = NULL;
    return((int)*optstr);
}

err(a0, e, c)
char *a0; /* argv[0]: the program name */
int e;
char c;
{
    if (!*errors[0])
    {
        strcpy(errors[0],
            "%s: illegal option - %c\n");
        strcpy(errors[1],
            "%s: option needs argument - %c\n");
    }
    if (opterr)
        fprintf(stderr, errors[e], a0, c);
}
```

CUT

The cut utility program allows the user to extract columnar fields from one or more files. It concatenates the fields together horizontally separating them by spaces or delimiters, and appends a new-line character after each line.

It is invoked in one of the following manners:

```
cut -clist [file-list]
cut -flist [-dchar] [-s] [file-list]
```

When used in the first manner, the list provides a set of character positions defining the columns.

When used in the second manner, the list provides a set of fields delimited by single characters.

The options are interpreted as follows:

-clist The list following -c specifies character positions (base one). For example, -c1-8,17-24,32 specifies character positions 1 thru 8, 17 thru 24, and 32. Either the -c or the -f option must be specified.

-flist The list following -f specifies field numbers (base one). These fields are separated by character delimiters, as specified by -d, or the tab character by default. For example, -c1-8,17-24,32 specifies field numbers 1 thru 8, 17 thru 24, and 32. Either the -c or the -f option must be specified.

-dchar The character following -d is the field delimiter character for the -f option. The default delimiter character is the tab character.

-s Suppresses lines with no delimiter characters with the -f option.

```
/*
 * This is a clone of the UNIX cut utility,
 *    except that the list of numbers
 *    does not have to be ascending.
 *
 * John Weald
 */

#include <stdio.h>
#include <ctype.h>

#define MAXLINE 1024 /* The max. length of a line */

main(argc, argv)
int argc;
char *argv[];
{
    extern char *optarg;
    extern int optind;

    FILE *fp; /* All the input files or stdin */
    char buf[MAXLINE]; /* The input buffer */
    char fs; /* The field separator */
    int c; /* The command line option */
    int cflag; /* True if -c on command line */
    int err; /* True if error in command line */
    int fflag; /* True if -f on command line */
    int suppress; /* Suppress lines with no delimiter */
    /* The field markers. True if this field */
    /* is to be cut, False otherwise */
    static int fields[MAXLINE];

    for (err = 0; err < MAXLINE; err++)
    {
        fields[err] = 0;
        buf[err] = NULL;
    }

    err = fflag = cflag = suppress = 0;
```

```
    fs = '\t';
    while ((c = getopt(argc, argv, "f:d:c:s")) != EOF)
    {
        switch (c)
        {
        case 'f': /* By Field */
            list(fields, optarg);
            fflag++;
            if (cflag)
                err++;
            break;

        case 'c': /* By character */
            list(fields, optarg);
            /* Implied suppress */
            suppress++;
            cflag++;
            if (fflag)
                err++;
            break;

        case 'd': /* A new field separator */
            fs = *optarg;
            break;

        case 's': /* Suppress if no delimiter */
            suppress++;
            break;

        default:
            prusage();
        }
    }

    if (!cflag && !fflag)
    {
        fprintf(stderr,
            "cut: Must have one of -f or -c\n");
        err++;
    }
    if (err)
        prusage();

    /* Loop on all the files. */
    do
    {
        if (optind == argc)
            fp = stdin;
        else
        {
            if (!(fp = fopen(argv[optind], "r")))
            {
                fprintf(stderr,
                    "cut: Failed to open file %s\n",
                    argv[optind]);
                exit(1);
            }
        }

        /* Loop on all lines in the file. */
        while (fgets(buf, 1024, fp))
            cut(buf, fields, fs, suppress, cflag);
        (void)fclose(fp);
    }
    while (++optind < argc);

    exit(0);
}

/*
 * Cut the line.
 * This handles both character and field cutting.
 * For characters f array gives character positions.
 * For fields it gives the field number.
 * It must be indexed by either the
 *    character number or the field number.
 */
cut(in, f, fs, sup, c_or_f)
register char *in; /* The input line */
```

```
    int f[]; /* The field cutting flags */
    char fs; /* The field seperator */
    int sup; /* Suppress lines with no-delimiter?   */
    int c_or_f; /* Cut by char. (true), or field (false)*/
    {
        char *instart; /* To print lines with no delimit-
ers */

        char obuf[MAXLINE]; /* Output buffer */
        register char *optr;
        register int fld; /* The field count */
        register int i; /* Character count */

        instart = in;
        optr = obuf;
        for (i = 0; i < MAXLINE; i++)
            obuf[i] = NULL;
        for (fld = 0, i = 0; i < MAXLINE; i++)
        {
            if (*in == '\n')
            {
                /* End of the line */

                *optr = '\0';
                /* Anything to cut? */
                if (optr != obuf)
                {
                    /* Get rid of trailing separator */
                    if (*(optr - 1) == fs)
                        *(optr - 1) = 0;
                    puts(obuf);
                    return;
                }
                if (!sup)
                    printf(instart);
                return;
            }

            if (f[c_or_f ? i : fld])
                *optr++ = *in;

            /* End of field? */
            if (*in++ == fs)
                fld++;
        }

        fprintf(stderr,
            "cut: Line too long, maximum length is %d\n",
            MAXLINE);
        exit(1);
    }


    /*
     * Parse the list argument. The format is:
     *    n,n
     * where n is either a number or a range of
     * numbers in the format
     *    m-l
     * m or l may be absent, indicating the start
     * or end of the lines respectivly.
     * Numbers must be in increasing order for m-l
     * format, but not for n,n.
     * Field numbers start at 1, but index into
     * fields array starts at 0.
     *
     */
    list (f, l)
    int f[]; /* The fields */
    char *l; /* The list */
    {
        int i;
        int high; /* the low and high numbers in a m-l
pair*/
        int low;
        int range; /* True if m-l format */

        low = 1;
        high = range = i = 0;

        while (1)
        {
```

```
            switch (*l)
            {
            case '\0':
                /* Is it m-<nothing>EOL? */
                if (range)
                {
                    /* Select rest of fields */
                    for (i = low - 1; i < high; i++)
                        f[i] = 1;
                }
                else
                    f[low-1] = 1;
                return;

            case ',':
                l++;
                if (!range)
                {
                    f[low-1] = 1;
                    range = 0;
                    low = 1;
                }
                else
                {
                    if (isdigit((int)*l))
                    {
                        low = atoi(l);
                        /* Skip the digits */
                        while (isdigit((int) *l))
                            l++;
                    }
                    range = 0;
                }
                break;

            case '-':
                l++;
                range++;
                /* Is it m-<nothing> */
                if (isdigit((int)*l))
                {
                    high = atoi(l);
                    /* Skip the digits */
                    while (isdigit((int) *l))
                        l++;
                }
                else
                    high = MAXLINE;

                /* Is the range the correct way around? */
                if (low > high)
                {
                    fprintf(stderr,
                        "cut: Bad c/f l)at: %d > %d\n",
                        low, high);
                    exit(1);
                }

                /* Set the field flags for the range */
                for (i = low - 1; i < high; i++)
                    f[i] = 1;
                break;

            default:
                /* either a number or an error */
                if (!isdigit((int)*l))
                {
                    fprintf(stderr,
                        "cut: Bad c/f list at %s\n", l);
                    exit(1);
                }

                if (!(low = atoi(l)))
                {
                    fprintf(stderr,
                        "cut: Fields start at 1 not 0\n");
                    exit(1);
                }
                /* Skip the digits */
                while (isdigit((int) *l))
                    l++;
```

```
                break;
            }
        }
    }


    prusage()
    {
        fprintf(stderr, "Usage: cut [-d<delimiter>] [-
s]");
        fprintf(stderr, "  -c<list>|-f<list> [file-
list]\n");
        exit(1);
    }
```

PASTE

The paste utility program allows the user to merge corresponding lines of several files or to merge subsequent lines of the same file.

It is invoked in one of the following manners:

```
paste [file-list]
paste -dlist [file-list]
paste -s [-dlist] [file-list]
```

In the first two manners, paste concatenates corresponding lines of the input files. It treats each file as columns of a table and concatenates them together horizontally.

In the last manner, paste combines subsequent lines of each input file.

In all cases, lines are separated by the tab character, or with characters from an optionally-specified list.

The options are interpreted as follows:

-dlist This option allows the specification of the field separation character list, which is a tab character by default. The list is used circularly. In parallel file merging (not with the -s option), the lines from the last file are always terminated with a new-line character, not from the list. The list may contain the special escape sequences: \n (new-line), \t (tab), \\ (backslash), and \0 (empty siring, not a null character).

-s This option may be used to merge subsequent lines from the same file rather than corresponding ones from each input file. Regardless of the contents of or absence of a list, the very last character of the file is forced to be a new-line.

- This option may be used in place of any file name, to read from the standard input.

```
/*
 * This is a clone of the UNIX paste utility.
```

```
 *
 * John Weald
 */

#include <stdio.h>

#define MAXLINE 1024 /* Max. allowed line length */
#define PLUSLINE 1025
#define MAXFILES 12 /* Max. number of input files */

main(argc, argv)
int argc;
char *argv[];
{
    extern int optind;
    extern char *optarg;

    char conchars[MAXFILES]; /* The concatenation
characters */
    int c; /* For getopt () */
    int nconchars; /* The number of conchars[] */
    int serial; /* True if type paste "-s" */

    conchars[0] = '\t';

    nconchars = 1;
    serial = 0;

    while ((c = getopt(argc, argv, "sd:")) != EOF)
    {
        switch (c)
        {
        case 's': /* Concatenate the same file serially
*/
            serial++;
            break;

        case 'd': /* Use other than a single tab */
            nconchars = setconcat(conchars, optarg);
            break;

        default: /* Does not return */
            prusage();
        }
    }

    if (serial)
        spaste(&argv[optind], conchars, nconchars);
    else
        paste(&argv[optind], conchars, nconchars);
    exit(0);
}

/*
 * paste()
 *
 * Do the actual paste.
 */
paste(files, con, ncons)
char *files[]; /* Null terminated list of input files */
char con[]; /* The concatenation characters */
char ncons; /* The number of above */
{
    FILE *fps[MAXFILES]; /* One for each open file */
    char c; /* The current concatenation char */
    char ibuf[PLUSLINE]; /* The input buffer */
    char obuf[MAXLINE]; /* The output buffer */
    int allfiles;
    int f; /* Number of files opened */
    int i;
    int inc; /* True if concatenation character == '\0'
*/
    int ocount; /* Output buffer character count */
    register char *iptr;
    register char *optr;

    iptr = ibuf;
    optr = obuf;

    /*
```

```c
         * Open all the input files, any filename of '-'
means
         * the standard input. No file name means standard
input.
         */
        for (f = 0; files[f]; f++)
        {
            if (*files[f] == '-')
                fps[f] = stdin;
            else
            {
                if (!(fps[f] = fopen(files[f], "r")))
                {
                    fprintf(stderr,
                        "paste: Failed to open file %s\n",
files[f]);
                    exit(1);
                }
            }
            if (f >= MAXFILES)
            {
                fprintf(stderr,
                    "paste: Too many files. Maximum
allowed is %d\n",
                    MAXFILES);
                exit(1);
            }
        }
        if (!files[0])
        {
            fps[0] = stdin;
            f++;
        }

        /* Read all lines until no more lines in any file.
*/
        allfiles = f;
        while (f)
        {
            optr = obuf;
            ocount = 0;
            /* Join lines from all files.
             * The concatenation character may be '\0'
which
             * means no character. The variable inc is an
indication
             * of the concatenation character being '\0',
we need to
             * if there is a concatenation character to
move up the
             * output buffer.
             *
             * The concatenation characters are used in a
round robin
             * list.
             */

            for (inc = 0, i = 0; i < allfiles; i++)
            {
                iptr = ibuf;
                optr += inc;
                /* To save repeated evaluation */
                inc = !(c = con[i % ncons]);

                if (!fps[i])
                {
                    /* No more lines in this file. */
                    *optr = c;
                    continue;
                }
                if (!fgets(ibuf, sizeof(ibuf), fps[i]))
                {
                    /* Reached EOF - finished with the
file */
                    (void)fclose(fps[i]);
                    fps[i] = NULL;
                    *optr = c;
                    f--;
                    continue;
                }
                /*
                 * Replace newline with the concatenation
character.
                 * There is no need to look for end-of-
string since
                 * we know that
                 * a) if ibuf is full to the max, then we
will
                 *    overflow obuf before we hit the end of
ibuf.
                 * b) if ibuf is not full, then it must
contain a
                 *    a newline character, but may or may
not
                 *    fit into obuf.
                 */
                for ( ; *iptr != '\n'; ocount++)
                {
                    /* Need space for trailing null */
                    if (ocount >= sizeof(obuf) - 1)
                    {
                        fprintf(stderr,
                            "paste: Output line too long,
maximum is %d.\n",
                            MAXLINE);
                        exit(1);
                    }
                    *optr++ = *iptr++;
                }
                *optr = c;
            }
            if (f)
            {
                *optr = 0;
                puts(obuf);
            }
        }
    }

    /*
     * setconcat()
     *
     * Parse the concatenation characters and place them in
the array c.
     * Return the number of concatenation characters.
     *
     * Specials are:
     *   \n    - Newline
     *   \t    - Tab (default)
     *   \     - Backslash
     *   \0    - No concatenation character
     */
    static int
    setconcat(c, in)
    char *c;
    char *in;
    {
        int i; /* The number seen so far */

        for (i = 0; *in; in++, c++, i++)
        {
            if (i > MAXFILES)
            {
                fprintf(stderr,
            "paste: Too many concatenation characters, more than
%d\n",
                    MAXFILES);
                exit(1);
            }
            if (*in != '\\')
            {
                *c = *in;
                continue;
            }
            switch (*++in)
            {
            case 'n':
                *c = '\n';
                break;
```

```
            case 't':
                *c = '\t';
                break;
            case '0':
                *c = 0;
                break;
            default: /* Includes '\\' */
                *c = *in;
                break;
        }
    }
    return(1);
}


static
prusage()
{
    fprintf(stderr, "Usage: paste [-s] [-d<list>]
files\n");
    exit(1);
}


spaste(files, c, n)
char *filee[]; /* Null terminated list of input files */
char c[]; /* The concatentaion characters */
int n; /* The number of above */
{
    FILE *fp;
    int i;

    if (!files[0])
    {
        spfile(stdin, c, n);
        return;
    }
    for (i = 0; files[i]; i++)
    {
        if (*files[i] == '-')
            fp = stdin;
        else
        {
            if (!(fp = fopen(files[i], "r")))
            {
                fprintf(stderr,
                    "Failed to open file %s\n",
files[i]);
                exit(1);
            }
            spfile(fp, c, n);
            (void)fclose(fp);
        }
    }
}


/*
 * Do the actual paste of a stream.
 *
 * The method here is to read in the stream and replace
 * newline characters with concatentaion characters.
 * Output occurs after each chuck is parsed, or if
 * the concatenation character is the null separator.
 *
 * The stream is read in BUFSIZ chunks using fread.
 * The input buffer is one character larger than read,
 * so that it can be null terminated.
 *
 * When we read in each chunk we must check if it
 * needs to be joined to the previous one.
 */
static
spfile(fp, con, ncons)
FILE *fp; /* serially paste this stream */
char con[]; /* The concatentaion characters */
```

```
int ncons; /* The number of above */
{
    char *pstart; /* The start of the string */
    char *ptr; /* Walks down the stream */
    char buf[BUFSIZ+1]; /* To ensure null termination */
    char last; /* The very last character looked at */
    int join; /* Join this chunk to the next? */
    int k; /* Index into concatenation character array */

    int n; /* Number of bytes read with fread() */

    join = k = 0;
    while ((n = fread(buf, sizeof(char),
        sizeof(buf) - 1, fp)))
    {
        if (join)
        {
            /* Join with last chunk */
            putchar(con[k]);
            k = (k + 1) % ncons;
            join = 0;
        }
        /* Join to next chunk? */
        if (buf[n-1] == '\n')
        {
            join++;
            /* Ignore the newline */
            n--;
        }

        /* ensure null terminated buffer */
        buf[n] = '\0';

        /* walk thru this chunk */
        /* replace newlines with next concat. char. */
        for (pstart = ptr = buf; *ptr; ptr++)
        {
            if (*ptr == '\n')
            {
                *ptr = con[k];
                if (!con[k])
                {
                    fputs(pstart, stdout);
                    pstart = ptr + 1;
                }
                k = (k + 1) % ncons;
            }
        }
        fputs(pstart, stdout);
        last = *(ptr - 1);
    }

    /*
     * Check for the newline as concatenation char.
     */
    if (last != '\n')
        putchar('\n');
}

EOF
```

PARDON IF I INTERRUPT

I have to relate an incident that happened to me some years back. It was about 4 years ago. I had bought a modem. It was a birthday present to myself. I was very enthusiastic about getting one. Around that time modems were the thing to have. Everyone had one. I knew many computerists whose only existence was to get on bulletin boards. And this became my goal.

When I got my modem, I decided that I too would go out and join the world of bulletin board systems (BBS). The first one I got on had a service that would list all the other BBS's in my area. What luck! I was in gravy. I downloaded the information to disk and later dumped it for a hard copy. So happy was I with my find that I even distributed to friends who wanted to get on the BBS's. I started to dial up different BBSs. I learned how to get on them. I found out how to use menus. I discovered how to get and leave messages. I was becoming a regular BBSphile.

The boards varied quite a bit. Some were information systems. I could learn many interesting things, like when and where ham fests were held, how to build a better modem and where the best local computer bargains were. Other systems were chat lines. People left messages for other people. One system was a 'joke board.' Its users left jokes. I even found one system that was devoted to dating. The theory was that members of the opposite sex could be met.

I got on one interesting system. After being a guest several times, the system asked me if I wanted to become a permanent user. The only cost was to register my name and some vital statistics. A few days later when I logged in it informed that I was now a member of the BBS and had rights and privileges far beyond what I had before. Entering 'HELP' would list all the features that were available. So I entered, 'HELP' and received a listing of what I could now do.

Among the commands was one marked, "SYSTEM." That sounded intriguing. My past experience indicated that this could open up even more doors. I would perhaps be able to get into the system and move through it freely.

So, I entered SYSTEM and immediately received the prompt "A:". Now this really did look familiar. It reminded of the system I had at work. I was sure now, I was in familiar territory. I entered, "DIR". Immediately a directory of a disk started to spew out on my screen, and then it stopped mid way in the listing. I typed on the keys and nothing I happened. I was latched up.

Then it started to print something like this:

Ha ha! I am the daemon of the computer. You have released me. For that I thank you. I am now in control. I can do anything I want.

Well, I was mildly amused. Obviously I don't believe in daemons. It continued.

First thing I will do is destroy all the hard disk information. I am now erasing all the files. One by one they are disappearing.

At this point I knew that no such thing could happen. No one would be silly enough to leave something like that in the system, unless another user had deposited it on the disk as a cruel joke. I immediately started to entering CONTROL-C, which on my work computer was always the means to stop a runaway program. No luck it continued.

I am now destroying the ROMs. Now I am now destroying the RAM memory.

I started to CONTROL-anything. I tried every thing. Finally in desperation I disconnected the modem line.

I thought about it. I knew it was a joke. But my curiosity got the better of me. I dialed the BBS again. I was greeted

as usual. It informed me of my status and then it said, "Ha! ha! Did I scare you?"

I replied, "No, but I was taken by surprise."

The sysop explained his little joke. He had put it on his machine for fun. I asked how he avoided being stopped by the CONTROL-C. Simple, he said, he had set a trap for the interrupt.

That brings us to this months topic, INTERRUPTS, SIGNALS and INTERCEPT TRAPS.

An interrupt is something that diverts the programs normal flow to some sequence of instructions that service the interrupt. It may be merely to acknowledge the interrupt or it may be something that requires immediate service. Interrupts are often used when some outside event requires attention. For example, a machine running under microprocessor control may be carrying out some prescribed task. An operator of the machine may place his hand in the machine endangering his safety. A sensor would pick this up and send an interrupt to the microprocessor. The program would immediately service the interrupt and stop the machine. This is an example of a hardware interrupt.

Signals are software interrupts. When you sit at the OS-9 keyboard and enter a CNTROL-E, your program will stop. What you have actually done is send it a signal. The signal was generated at the keyboard and told the process to stop.

The signal can therefore be thought of as asynchronous control device for communica-

tions between processes. It is like the hardware interrupt from before. How the program responds is determined by the software.

In OS-9 the signal is a 1-byte numeric value. It is sent from one process to another. The meaning of the signal has been predefined for some values. The rest of the values are user definable. Here is list of them:

0 - Kill
1 - Wakeup
2 - Keyboard terminate
3 - Keyboard interrupt
4 - Window change
5-255 User definable

Signals sent to a process are saved in the process descriptor. If the process is sleeping or waiting, it is moved to the active state. On the next time slice, the signal is processed.

How it is processed, depends on the program. An interrupt handler can be built into the process to take care of signals. It easy to find an example. When in BASIC09 and a program encounters an error, it goes into DEBUG mode. Using a CONTROL-E will place it back in the SYSTEM MODE. It does not cause it to abort and go back to the OS-9 shell.

If a process has not taken steps to handle signals, then it will be terminated immediately. Most programs that are written fall into this category. Try a CONTROL-C or CONTROL-E from the keyboard and your process will stop. In most cases this is a desirable outcome. It is rather upsetting to start a process and then find it can not be easily stopped. So by doing nothing, you have a program that is easily haltable.

A process can take measures against an incoming signal. The process sets a trap. When a signal is sent to the process, it is passed to the intercept routine. What the routine does with the signal code it receives is up to it. The routine always ends with a RTI. This insures that normal execution of the process. It will return to whatever it was doing before the interrupt.

There are some exceptions to what I have just spelled out. For example, if the signal sent is 0 which is to kill the process, it will die. If this signal were trapable, it would mean that any process could runaway with no stopping it. When the SHELL command KILL is executed, it uses this to halt the process.

KILL is a command that actually part of the SHELL. It syntax is

KILL processID

Unless you are Super User, you can only terminate you own processes. A process will not die if it has any pending I/O's. If you ever use KILL on a process and it still exists, it means it is has something going on. Most likely, it is reading or writing data somewhere.

Another signal that gets a different treatment is the wakeup signal. It wakes a sleeping process. However, the trap routine is never entered.

One other thing about signals. If a process is sleeping and it receives an a signal, the signal will be pending until the process awakes and gets its next time slice. If another signal is sent to the process, an error will be sent to sender. The sender could then go to

sleep for a few ticks and try again. (A call to F$Sleep saves CPU time.)

A call F$Send will send a signal to any process. All that is necessary is to load the A register with process ID and the B register with the signal. Most cases the signal go through. If it does not, an error will return. Otherwise, the process will receive the signal. If it is in a Sleep or Wait state, it will be activated on its next time slice.

A receiving process should set up an intercept trap using F$Icpt. If it does not do this it will terminate with the incoming signal. The X register is loaded with address of the intercept routine and the U register with the routine's memory area. When a signal is sent to the process, execution is transferred to the routine. The U register will contain a memory area for its use and the B register will have the signal code. The routine must be terminated with a RTI.

Anything can be done within the intercept routine. The simplest method is to ignore it. This is what the SYSGO process does. Its routine is only a simple RTI.

The routine may be more elaborate. If a signal code of 2 or 3 comes in, instead of terminating right away, some housekeeping can be done. Buffers can be emptied to files. Files closed. Temporary files renamed. Whatever it takes to successfully exit your process can be done. If you don't want to exit, take care of it any way you wish.

To better illustrate how to set an intercept trap for incoming signals I have included a program this month called DAEMON. It is my version of

the program that I described earlier. You leave this lying around your commands directory for some unsuspecting soul to find. By its very name, it should get attention immediately. Once they enter it, it will announce its arrival and then obstinately refuse to leave. Some of the quips it throws out are:

You can't get rid of me! I will haunt you forever. You fool! OS9:... just kidding!

I'll let you read the rest.

I set up an intercept trap called catchit(). I used the C Language call intercept() that comes with the Microware compiler. This call is some similar to F$Icpt. The location of the catchit is passed. This sets up where to pass program execution whenever a signal comes in.

The C call to intercept() as I said is similar to F$Icpt. In reality, the call passes a pointer of where the C routine is for handling the signal. Then it sets up a routine called receiver which is the true signal trap. When entered it pushes the signal code onto the stack and does a JSR to the C routine. In the case of my program it is catchit().

The program catchit assumes it will be passed one value. It is an integer. The value of which is the signal code. I have anticipated a 2 for Keyboard Abort and 3 for Keyboard Interrupt. Just in case something else comes in, I have added a catch all phrase.

The KILL command can still stop DAEMON. So also replying 'Get Lost' will stop it. I left the 'Get Lost' in as a safe guard to get out and as a challenge for someone else to escape the DAEMON.

Other interesting items about daemon.c is that I created a routine to return a psuedo random number from 0 to 9. My C libraries do not have any math functions. So I am left create one. The routine is called next(). What it does is to take the last number returned, add to it the length of the last line of input and add the ASCII value of the first character. This it divides by 10 and returns the remainder. OK, it's cheap and dirty, but it works

I stuck strictly with the standard I/O calls — writeln and readln. I had originally used printf() and gets(), but had problems when using it with my intercept trap. Rather then delve into the project in great depths, I took the standard way out. I programmed around it.

I let you to try this one. I might be fun, unless you are the one who is stuck with trying to get out of it.

Have a good month. See ya next time.

```
0000 /*  ****************
0001
0002       Name: Daemon
0003       Date: 28-Feb-88
0004       By:   Ron Voigts
0005       To compiler: cc
daemon
0006
0007       ****************
0008
0009       Usage:
0010       Leave in commands
directory
0011       for some unsuspect-
ing user.
0012
0013       To stop it, enter
'Get Lost'
0014       or use KILL command.
0015
0016       ****************
0017
0018       Verison 1.0    Origi-
nal  ROV
0019
0020       ****************
*/
0021
0022 #include <ctype.h>
```

```
0023 #define MAXLINE 80                                 0083 catchit( n )
0024                                                      0084 int n;
0025 main()                                              0085 {
0026 {                                                    0086       if ( n==2 )
0027       /* Declarations */                             0087           printit("Keyboard Abort doesn't
0028       int catchit();                            work on me.");
0029       char s[MAXLINE];                               0088       else if (n==3)
0030       int i;                                         0089           printit("Keyboard Interrupt won't
0031                                                  work");
0032       /* Responses of daemon */                      0090       else
0033       static char *r[] = {                           0091           printit("Can't get rid of me, eh");
0034          "You can't get rid of me!",                 0092 }
0035          "I am the Deamon of the Computer.",         0093
0036          "I will haunt you forever.",                0094 /* Routine to print a string line */
0037          "Now what are you going to do?",            0095 printit( s )
0038          "Ha-ha.  I am in control.",                 0096 char *s;
0039          "You fool!",                                0097 {
0040          "Are you sorry you invoked me?",            0098       int i;
0041          "I might stay here forever.",               0099       i = strlen( s );
0042          "OS9: .... just kidding!",                  0100       writeln(1,s,i);
0043          "Well, what are we going to do?",           0101       writeln(1,"\n",1);
0044       };                                             0102 }
0045                                                      0103
0046       /* Set the trap */                             0104 /* Cheap random number generator */
0047       intercept( catchit );                          0105 next( s )
0048                                                      0106 char *s;
0049       /* Opening message */                          0107 {
0050       printit("I am the daemon of the com-           0108       static int i;
puter.");                                                 0109       i+=strlen( s )+s[0];
0051       printit("You have just released me from        0110       return( i % 10 );
the disk.");                                              0111 }
0052       printit("For this I thank you.   But now       0112
I must");                                                 0113 /* Prints the 'Well?' prompt */
0053       printit("do what ever damage I can.            0114 prompt( s )
Just wait");                                              0115 char *s;
0054       printit("and see what I can do.");             0116 {
0055                                                      0117       printit("Well? ");
0056       /* Main loop */                                0118       getit( s );
0057       for (;;) {                                     0119       writeln(1,"\n",1);
0058                                                      0120 }
0059          /* Get input from user */                   0121
0060          prompt( s );                                0122 /* Gets an input line */
0061                                                      0123 getit( s )
0062                                                      0124 char *s;
0063          /* Put it into uppercase */                 0125 {
0064          toupperc( s );                              0126       int i=0;
0065                                                      0127       readln( 0, s, MAXLINE );
0066          /* This will end the process */             0128       while( s[i]!='\n' )
0067          if ( strcmp( s, "GET LOST" )==0 )           0129           i++;
{                                                         0130       s[i]='\0';
0068              printit("You got me, you dirty          0131 }
dog!");                                                   0132
0069              exit( 0 );                              0133 /* Coverts a string to uppercase */
0070          }                                           0134 toupperc( s )
0071                                                      0135 char *s;
0072          /* Get pseudo random number */              0136 {
0073          i = next( s );                              0137       register int i;
0074                                                      0138       for ( i=0; i<strlen(s); i++ )
0075          /* Print the daemon's response */           0139           s[i]=toupper( s[i] );
0076          printit( r[i] );                            0140 }
0077                                                      0141
0078       }                                              0142
0079
0080 }
0081
0082 /* Program to handle incoming signals */
```

EOF

FOR THOSE WHO NEED TO KNOW   **68 MICRO**™
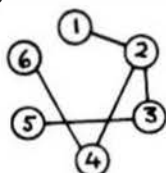   **JOURNAL**™

Logically Speaking

The Mathematical Design of Digital Control Circuits

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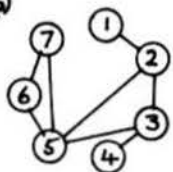Canada V2S 1E2

Solutions to TEST EIGHT

1.

2. There are several possible alternatives, so only one solution is given as being typical of the method of introducing dummy states to ensure a successful Gray-coding.

3. For the purposes of this design we'll call the push-button X1, the revolution-indicating micro-switch X2, the table-motor Z1, the soap control valve Z2 and the water control valve Z3. The small diagram at the top-right below shows how the switch signals, with the operating-arm dropping into a detent at the end of each revolution. This is also the starting-point of the cycle.

In order to get you started on the flow-table, let me explain the first few steps. We begin at address 00.1, with all Zs OFF, then when button X1 is pressed nothing happens to the Zs but we move to address 10.2. On releasing X1, we move to address 00.2, where all Zs become energised, and, of course, as soon as the table starts to rotate the lever on X2 operates, moving us to address 01.3, and maintaining us in this state for the duration of the first revolution. At this point in the revolution, X2 drops down into the detent, de-activates, and moves us to address 00.4, still maintaining all Zs ON.

And so on through successive revolutions of the table, until the end of the third rev. when the soap is shut OFF, with a phi at the elbow. For the final rev we move back to our starting row 1, so that when X2 drops out for the last time we pop back to address 00.0. End of operation!

The intermediate decoding-table shows that the operating conditions for both table-motor and water-valve are the same, so only five decodings will be necessary.

$$Y1 = X2'.y1 + X2.y2.y3' + y1.y3$$
$$= y1(X2' + y3) + X2.y2.y3'$$
$$Y2 = y2.y3' + X2'.y2 + X2.y1'.y3$$
$$= y2(X2' + y3') + X2.y1'.y3$$
$$Y3 = X1 + X2.y3 + X2'.y1.y2 + y1'.y2'.y3$$
$$= X1 + y3(X2 + y1'.y2') + X2'.y1.y2$$
$$Z1, Z3 = X1'.y3 + X2 + y1 + y2$$
$$Z2 = X1'.y3 + y2$$

I'll leave the drawing of the circuit-diagram to you. as I've already done the factoring to help you along.

How did you make out with that little test? I hope you're getting more used to this stuff now. but it'll get re-inforced a bit more with what's to come. Here we are at

Mile 10 - heading for Mile 11.

TIMERS AND SELF-ACTUATED MICRO-SWITCHES

We'll assume that our last customer was so pleased with the circuit we designed for his car-wash that he's recommended an associate to us. who now wishes us to design the following control system.

DESIGNING A PRESS CONTROL

We have a press. and a push-button X1. such that when X1 is operated the press will lower and operate a micro-switch X2 at the bottom of its stroke. The press must then squeeze for 10 seconds before returning to the top of its stroke. The squeezing operation is to be a "one-shot" affair. That is. only one operation will be performed whether X1 is maintained or not. X1 MUST be released. and re-operated at the start of the cycle. in order to initiate another sequence. Re-operating X1 DURING the cycle will be treated as though it were never released in the first place.

Diagram 45

This is a new kind of machine for us to design. so we'll cover the flow-table of Diagram 45 in some detail. Z1 will be the press-actuator. and Z2 will be the energising-coil (or motor) for the timer. Because the squeeze-period is a stable state (at least for 10 seconds). and the timer is going to be responsible for causing our machine to leave this state. its contacts will be classed as a primary control. appearing as T in our column-header. Think of it. if you like. as an automatic-X3. which saves the operator from timing the 10-seconds himself and then having to press X3 at the end of this time. Just as the micro-switch X2 is also an automatic device. which saves the operator from having to watch the

press and then hitting X2 himself at the bottom of the stroke. The press COULD, of course, take quite a few seconds to reach the bottom of its stroke!

To summarise, delayed devices of this kind will have their coils, motors, etc, classed as Zs, but their contacts as Xs (primary controls).

Now for the construction of the flow-table. We'll commence with a table which is blank except for a few rows numbered, let's say, from 1 to 6, and a heading X1.X2.T together with one column, headed 000 (ie, all Xs OFF). Our starting-point is, by convention, address 000.1, where we hold the circuit action stable with a 1 in Box-A, and 00 in Box-C to keep both the press and the timer de-energised. There's only one way to get out of this stable state, and that is to press X1, which is the only control directly available to us. This forces us to create a column with heading 100, and moves us down to address 100.2, with a phi at the elbow for the press energisation. Keeping in mind that the cycle has to be identical even if X1 is released, we'll duplicate the codings of address 100.2 in address 000.2.

At this stage, the press is on its way down to work on the article beneath it, so nothing further is going to happen till it reaches the bottom of its stroke and hits X2. When it does so, our timer must start timing, so our next entries are fairly straightforward. We create two new column headings (one assuming that X1 is still operated and a parallel one assuming it's been released), both of which indicate that X2 has now been actuated. The circuit-action moves through both columns in parallel, with identical codings in each, to addresses 010.3 and 110.3, at which point we now energise the timer (keeping the press energised), with our usual phi on the elbows.

We're going to sit here in a stable state for 10 seconds until timer-T times out and signals the end of the squeeze-period. Again we create a new, parallel set of columns, and move the circuit-action into addresses 011.4 and 111.4, where we de-energise the press, with a phi on the elbow. Note that we don't de-energise the timer as well, as we're VERY conscious of the fact that only one thing should happen at a time in our circuits. For instance, if we DID de-energise timer-T as we move from stable-state 3 (addresses 010.3 and 110.3) to unstable-state 4 (addresses 011.3 and 111.3), there would be a critical race between the contacts of timer-T trying to drop out (and move us back into stable-state 3) and the relay-action trying to move us down into stable-state 4 (addresses 011.4 and 111.4), and there'd be no telling which would win!

Anyway, by keeping timer-T energised, we'll definitely end up in either 011.4 or 111.4, with the press freshly de-energised and ready to start on its upward journey back to the start of the cycle. The next thing to occur is that on its way up it moves off micro-switch X2, forcing us to open up two more parallel columns, headed 001 and 101. Now, as this represents the last stage of the cycle, we'll examine separately the paths of action through the flow-table, (a) with X1 NOT operated and (b) with X1 still held down.

We're therefore in the process of moving from address 011.4 back to the start of the cycle. Obviously, we must get back to row 1, so we enter a 1 in Box-A of both 001.4 and 001.1. Further, we'll keep timer-T energised during this vertical transition (no phi on the elbow) to ensure that its contacts, in resetting, do not enter into a critical race with those of the relay responsible for moving us vertically through the flow-table. When we reach row 1, AND ONLY THEN, will we insert a 0-entry in Box-C to de-energise the timer. This, of course, opens its contacts (which closed at the end of the 10-second timing period) and very nicely moves us back to our starting-address 000.1. Are you with me so far? Good!

So let's construct the parallel path, commencing at address 111.4, then we're all done with the flow-table. Again we wish to move up to row 1 in order to get back to the beginning of the cycle, so, just as before, we enter 1s in Box-A of addresses 101.4 and 101.1. However, this time we do NOT de-energise timer=T when we arrive at address 101.1. This would surely shoot us to column 100, where we'd find an instruction '2' in Box-A, and the machine would keep on repeating its cycle until such time as we finally decide to release X1. We therefore decide to keep T energised in address 101.1, where we'll sit indefinitely until we release X1, which would slide us into the adjacent address 001.1, and we'd terminate as we did for case (a).

AN ALTERNATIVE ENDING TO OUR TABLE

Before we carry on with the actual design, let's look at an alternative way to get back to row 1 from row 4. Let's start again from address 011.4, where X2 is due to release itself because the press has

just been de-energised. When it does we'll move into address 001.4, just as before, but this time we'll enter a 4 in Box-A to hold us in row 4, and de-energise timer-T at the same time. There's no question of a critical race now, as no relay action is taking place, so timer-T will release its contact and move us into address 000.4. Box-C of this address should have a 00 entry to keep both the press and timer de-energised, and Box-A a 1 to return us to the starting-point.

Repeating this action from address 111.4, we'll move into 101.4 with a 4 in Box-A and 00 in Box-C to de-energise timer-T. Again timer-T will release its contact, moving us now into address 100.4, where we'll enter 4 in Box-A. Do you see why we shouldn't enter a 1 here? We'll now remain stuck in this location until we release X1, and thereby move to address 000.4 and so back to the start.

When we come to merging rows, however, you'll see why this is not as good as our first try. Therefore, when we plan our flow-table action, we'll try to keep in mind the fact that some paths will lend themselves to a better merge-pattern than others.

CHANGES IN SPECIFICATIONS

It could have been the case that our client wished the machine-sequence to keep repeating if the main operating-button X1 were held depressed, and to complete the current cycle if X1 were released at any time during the cycle. We can see from our flow-table that, although this sounds quite a drastic change in specs from those of our machine here, the only difference would be an entry of 00 in Box-C of address 101.1, instead of our current 01 entry. On the other hand, it sometimes happens that an apparently small change in the specs can lead to the most drastic upheaval in the flow-table, even to the extent of having to scrap the one just constructed and to start all over again.

CONTINUING WITH THE DESIGN

Well, having got this far, we may as well go on to our merger-diagram, shown in Diagram 46a. You should try for yourself the merger-diagram for the alternative "return-to-start" path which we considered above.

$x_1 x_2 T$	0 000		8 100		4 010		12 110		6 011		14 111		2 001		10 101	
1	1	0	2	8	1	4	1	12	2	6	2	14	1	2	1	10
0	00		φ0		11		11		φ1		φ1		00		01	
2	2	1	2	9	1	5	1	13	2	7	2	15	1	3	1	11
1	10		10		1φ		1φ		01		01		01		01	

(a)   Diagram 46   (b)

With experience, our original flow-table would normally be constructed with less rows than those of Diagram 45, but the action is clearer and a further demonstration of merging is possible with the 4-row flow-table which we developed. The merger-diagram shows that we have two choices of merging rows, both of which would result in a 2-row flow-table. That is, we can merge rows 2, 3 and 4, and leave row 1 as is, OR we can merge rows 1 and 3, and rows 2 and 4. For the purposes of our example we'll try the latter choice, and call rows 1 and 3 our new row 1, and 2 and 4 our new row 2 (see 46b). Both rows are completely filled with entries, but it should be noted that if we'd taken the first choice four addresses of row 1 would be phi-states. Maybe you should try doing this, and see if the resultant circuit is any simpler than the one we're going to do here.

It won't be necessary to do a state-diagram for a 2-row flow-table. One relay will do the job, and we have no choice therefore but to code row 1 with 0 and row 2 with 1, and to allocate to X1, X2 and T the bit-positions 8, 4 and 2 respectively. All the Box-Bs are coded in red decimal figures, according to the sum of their co-ordinates, and the three required outputs (Press, Timer and Y1) are then decoded to give us the Boolean expressions from which we can quite easily construct the circuit-diagram.

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
(a) Z_1	1			1	1	φ		φ	1				1	1	φ	
Z_2			1	1	φ	1	1				1	1	1	φ	1	1
Y_1		1					1	1	1	1					1	1

(b)

Z_1

	1	4	5	9	12	13	X_1	X_2	T	y_1	φ
x		4	8		4		φ	φ	0	1	
x			8				φ	1	φ	0	6,14 ✓
x	1			8	8		φ	1	0	φ	✓

$\text{Press} = T'y_1 + X_2 y_1' \quad \text{OR} \quad \text{Press} = T'y_1 + X_2 T'$
$$= T'(y_1 + X_2)$$

(c)

Z_2

	3	4	6	7	10	11	12	14	15	X_1	X_2	T	y_1	φ
x		4		8					4	φ	φ	1	1	
x	2	1			8		2	1		φ	1	φ	φ	5,13
x	1					4	1			1	φ	1	φ	

$\text{Timer} = Ty_1 + X_2 + X_1 T$
$$= X_2 + T(X_1 + y_1)$$

(d)

Y_1

	1	6	7	8	9	14	15	X_1	X_2	T	y_1
x				8				φ	0	0	1
x	1				8		1	φ	1	1	φ
x	1							1	0	0	φ

$Y_1 = X_2'T'y_1 + X_2 T + X_1 X_2' T$
Diagram 47 $= X_2 T + X_2'T'(X_1 + y_1)$

You'll notice in 47b (the decoding for the Press, Z1) that we're blocked in our attempt to eliminate variable y1 in row 2, the block being due to the presence of the phi minterms 6 and 14. We therefore try an auxiliary decoding-row 3, where we eliminate y1 first, and then continue normally. This gives us an alternative decoding for Z1, which, by good fortune, is capable of being factorised into a 3-literal expression instead of the original 4-literal one.

Again I'll leave you to draw the circuit-diagram, as I don't want you to get muscle-bound through lack of exercise! And while you're all loosened up, I'll give you just one test question for TEST NINE, because next time around .... guess what I have in store for you??? This will make you TRULY happy! I'm going to show you how to construct a very simple, VERY effective set of cards which will enable you to perform the decoding operation almost automatically. Naturally, I hope you'll make yourself a set!!

TEST NINE

1. Design the following

An alarm light is initially OFF. When X1 is operated, it remains OFF. Provided X2 is operated at any moment greater than 2 minutes, but less than 3 minutes, from the time X1 is released, the circuit will return to its initial state. However, if X2 is operated before the 2-minute period expires OR after the 3-minute period, the alarm light will come ON. This is where our design will stop, as we'll assume that once X2 is operated it stays operated, and that a Master-Reset (also not to be included in our design) will be required to kill everything and return it to its initial state.

... End of Mile 10

EOF

FOR THOSE WHO NEED TO KNOW   68 MICRO JOURNAL™

SOFTWARE USER NOTES

A Tutorial Series

By : Ronald W Anderson
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Ann Arbor, MI 48105

From Basic Assembler to HLL's

68000 Assembler

As you all might guess from last month, the bug has managed to bite me rather hard. This and probably some of the following columns will contain some words on 68000 Assembler. Now please don't expect expert advice from me on the specifics of the 68000. I am learning as I go here, but I've noticed from past attempts that the responses seem greatest and most positive when I write about present experiences and problems with programming. Last time I threw three programs at you. No sooner had I mailed off the material to '68' Micro Journal than I heard from Peter Stark. He told me that I had a serious bug in the two programs PLIST and UPCASE. Both have the same problem. I had assumed that A6 would be pointing at the user FCB area on starting the program. The SK*DOS manual indicates that this is not necessarily true. A6 points at the user FCB area only after an SK*DOS system call has been made. Any system function will do, and if you don't want to do anything first, SK*DOS has provided a do-nothing function called VPOINT that does nothing but set A6 to pointing at the FCB area. Therefore, my program PLIST should not have as its first instruction:

START MOVE.L A6,A0

Rather that instruction should be replaced with the following two lines:

START DC VPOINT
MOVE.L A6,A0

In addition, the function VPOINT must be added to the list of equates for the SK*DOS functions at the beginning of the program. VPOINT EQU $A000 is all it takes. The program UPCASE needs exactly the same fix. The reason that the programs worked, was that A6 pointed into the stack area where an FCB could be set up without causing disaster. When I tried to access one of the SK*DOS data locations with another program I had a problem and this turned out to be the solution.

Now rather than diving right in with some further examples of programs, maybe we ought to discuss the 68000 a little bit. You 6809 users will remember that the 6809 has two 8 bit accumulators called A and B, two 16 bit index registers X and Y, and two stack pointers U and S. The S stack pointer is used by the processor as the "system" stack pointer but the U or User stack pointer may be used by the programmer. In addition, the X and Y registers, though primarily designed as index registers, and the U register primarily designed as a stack pointer, could be used more or less interchangeably. Because the accumulators had names (A and B) these names could be used in instructions that referenced them. LDA #5 would transfer the value decimal 5 to accumulator A. The 6809 instruction set is rather regular, but not as regular as the 68000 for reasons that will be obvious shortly.

The designers of the 68XXX processors chose to provide a larger number of registers that could be used as accumulators. That is, all operations can be performed on any of the accumulators, which in the 68000 are called data registers. There are EIGHT of them, D0 through D7. All are 32 bit registers in all of the 68XXX processors. These registers are used for all operations.

In the 6809 you couldn't ADD #3, to a memory location. You first had to LDA LOCATION, then ADDA #3, then STA LOCATION. With the 68000, some operations may be performed directly into memory. ADD.B #3,(A0), for example will add 3 to the contents of the memory location whose address is in register A0. Of course all of the operators work on data in registers. In 6809 code you use LOCATION RMB 2 to reserve two bytes of memory for your program to use to store values. RMB means Reserve Memory Bytes. With the 68000 you use LOCATION DS 1. DS means Declare Storage. DS.B 1 will reserve one BYTE. DS.W 1 or just DS 1 will reserve one Word or two bytes. DS.L 1 will reserve a long word or 4 bytes. The 6809 assembler has instructions for declaring constants in a program too. FCB 8, for example means Form Constant Byte with a value of 8 decimal. FDB $1234 means Form Double Byte (HEX) 1234. Constants are declared in the 68000 assembler just about the same way as variable storage. DC $1234 declares a Constant $1234. DC.B and DC.L are used just as with DS above.

The 68000 assemblers have a further generalization that I found not dificult to use. Most of the instructions will work on operands of different sizes. ADD.B #3,D0 will add a byte value of 3 to the low order byte of the D0 register. The three higher order bytes won't be changed at all. Similarly ADD.W #3 will add 3 to the low order 16 bit value in D0, the upper 16 being ignored. Lastly, ADD.L #3,D0 will add 3 to the register, operating on all 32 bits. The suffixes .B, .W, and .L indicate Byte, Word, and Long, respectively. If no suffix is assigned by the programmer, the .W is assumed. Not all instructions work on all three data sizes. The only way to be sure is to consult the 68000 instruction set in one of the Motorola publications, or to try it and see if the Assembler balks.

In addition to the 8 data registers, the 68000 has 8 address registers A0 through A7. The system reserves A7 as its stack pointer. The user can use A0 through A6 in his programs. However, SK*DOS uses and alters some of the registers when you call one of its routines. You are perfectly free to use D0 through D3 and A0 through A3 for your user program with guarantees that SK*DOS will not change them. Some of the function calls to SK*DOS require that you pass information in one or more of the registers. PUTCH, the routine to output a character to the terminal, for example requires that you put the character in D4. To print a space you would MOVE.B #$20,D4 DC PUTCH.

When you call other SK DOS functions, you get information back in a register too. For example DC GETCH will return a character from the terminal in D5. When you go to print a text string terminated with $04, you point A5 at the first character and DC PSTRNG. Maybe a comparison here with the 6809 way would be helpful. 6809 FLEX would use LEAX MESSAGE,PC to get X to point at the message in a position independent assembler program. That is, X would be loaded with the address of the first byte of the MESSAGE string relative to the present program counter. Then you would JSR PSTRNG. PSTRNG would be defined by equating that label to an address somewhere early in the program. Now with the 68000 you use something very similar. LEA MESAGE(PC),A5 DC PSTRNG. While in the 6809 FLEX environment, you had to jump to the subroutine in FLEX to do the job, the 68000 SK*DOS works more like a software interrupt. Op codes never start with the byte $A0, and that is an illegal instruction. DC PCRLF, for example, causes the value $A034 to be placed in the next two bytes of the program. The 68000 traps the illegal instruction byte $A0, and the trap handler in SK*DOS interprets the next byte to decide which routine to jump to. Note that these values are all 16 bit or Word values so we don't need a suffix on DC. p Well, what good are the Address registers anyway? Like those X, Y, and U registers in the 6809, they can be used as index registers or stack pointers. LEA MESSAGE(PC),A0

points A0 at the first byte of MESSAGE. Now MOVE.B (A0),D4 would transfer that first byte of MESSAGE to D4. Then DC PUTCH would output it to the terminal. The notation (A0) means to use A0 as a pointer. That is, get the data via what is called register indirect addressing. It is just like LDA 0,X in the 6809 notation. Actually, if we were writing a simple section of code to output a string we would probably want A0 to advance and point at the next character in the string. We could do that by ADD.L #1,A0, or much more easily by using the post increment addressing mode. MOVE.B (A0)+,D4 would move the character and increment A0 all ready for the next character. In fact we could write the PSTRNG function (not worrying about carriage returns or linefeeds) simply now.

```
PSTRNG LEA
MESAGE(PC),A0
LOOP   MOVE.B (A0)+,D4
       CMP.B #$04,D4
       BEQ.S EXIT
       DC PUTCH
       BRA.S LOOP
EXIT   RTS
```

This code would first point A0 at the message, get the character in D4, incrementing A0 after the transfer. Then it compares the character with $04 which, you remember, ends the text string. If it is $04, the loop ends with a branch to EXIT which returns from the subroutine. If it is not $04, we go around the loop again and get another character and output it. The comparison tests are just like those of the 6809. BEQ means to branch if equal after the comparison of $04

to the character in D4. All branches may have the suffix .S for short if they are within the range of +127 to -128 bytes. Without the suffix, long branches are assumed. Some assemblers automatically select short branches if they are within range, others do not. Generally, you can assume that if a branch is to a label that can appear on the same screen as the branch when you edit the source file, you can use a short branch. If in doubt, use a long one and see if the code generated is less than 00007F for a forward branch or greater than FFFF80 for a backward one.

The instruction set, with that much explained, is not difficult to understand except for some special instructions that will be used by compiler writers more than by everyday Assembler programmers. There are ADD and SUB instructions and like ones that take into account the carry/borrow (simply called the Extend bit in the 68000). These instructions are ADDX and SUBX. They may be performed on byte, word, or long data. There is a CLR instruction that can clear a register or memory location. There are Rotate and Shift instructions, both left and right, and operating on bytes, words, and longs, both including and excluding the extend bit. There are signed and unsigned multiply an divide instructions, shift N bits instructions, limited to 8 places if the data is supplied in IMMEDIATE mode and any number if a register holds the shift count. ASR.W #7,D0 will shift the low order 16 bits 7 places to the right, ignoring the upper 16 bits. ASR.L D1,D0 will

shift the contents of D0 by a number of places specified by the contents of D1.

To be quite honest, when I write little assembler programs in 68K assembler I feel more as though I am working with a higher level language. In terms of listing size, some of my simple utilities have listings that are far shorter than the equivalent program in 6809 assembler. The code generated is about 40% more, however. That is partially explained by the fact that the 68000 addresses a much larger address range and that it operates on 4 byte operands as well. It just takes more bytes to code an instruction in 68K. I found the transition to be painless and quick. True there is a new syntax to learn but there is suddenly so much more available to work with! For some of the simple utilities I

have written, I have not had to declare any variables at all by using registers to store counters, pointers, and intermediate results.

The listings supplied this time should serve to illustrate some of the above points. There is a program to dump a page of memory in HEX and ASCII. MDUMP gets you into the program. It prompts for a command. The 68008 system that I have has memory from 000000 on up to 0BFFFF. You specify a page on which to start, for exmple by typing N005E. The program indicates 005E00 as the starting address for the dump and dumps that address through 005EFF, 16 bytes to a line with 16 hex values followed by their 16 ASCII equiva- lents. Non-printable codes are shown as periods. To dump the next page of

memory, you simply type F for Forward. To back up one you type B for Back. To go to a new area of memory you use N again. To exit either Q or E will return you to SK*DOS.

DDUMP is the same program except that it will dump a disk file one sector at a time. You type DDUMP followed by the filename and extension. It will dump sectors until you get to the end of the file and then exit to SK*DOS. If you have PAUSE enabled, it will stop after each sector.

FIND is a utility that was available for FLEX. I didn't translate it, but started from scratch. You type FIND and a FILENAME. .TXT is the default extension but you can supply any other. The program then prompts "INPUT SEARCH STRING". You supply the

string followed by CR and FIND prints out each line of the file that contains the search string, along with a line number.

After conferring with Peter Stark, I decided to add the help text to each of the utilities. Incidentally, I've made disparaging remarks before about endless CAT, DIR, COPY utilities published for FLEX. I still feel that way about them, but there certainly is no better way to learn assembler programming and a little about the operating system, than to write some simple utilities and make them work. After writing 16 to 35K object code programs for a living for a while, it is very refreshing to write 200 to 500 byte utilities in assembler.

```
list mdump93


      1                    *
      2                    * MEMORY DUMP PROGRAM
      3                    *
      4                    * COMMANDS:
      5                    * N XXXX NEXT PAGE TO BE DUMPED
      6                    * B BACK A PAGE
      7                    * F FORWARD A PAGE
      8                    *
      9                    * SK*DOS / 68K EQUATES FOR USER PROGRAMS
     10
     11         0000A029   GETCH  EQU    $A029          Get input character with echo (7 bits)
     12         0000A02D   GETNXT EQU    $A02D          Get next char from command line
     13         0000A031   TOUPPR EQU    $A031          Conv't char in D5 to Upper Case
     14         0000A02F   HEXIN  EQU    $A02F          Input hexadecimal number
     15         0000A03A   OUT2H  EQU    $A03A          Output 2 hex digits
     16         0000A03C   OUT8H  EQU    $A03C          Output 8 hex digits
     17         0000A034   PCRLF  EQU    $A034          Print CR/LF
     18         0000A036   PNSTRN EQU    $A036          Print string (Without CR/LF)
     19         0000A035   PSTRNG EQU    $A035          Print CR/LF and string
     20         0000A033   PUTCH  EQU    $A033          Output character
     21         0000A000   VPOINT EQU    $A000          Point to SK*DOS variable area
     22         0000A01E   WARMST EQU    $A01E          Warm start
     23
     24                    *
     25                    * COMMAND PROCESSING LOOP
     26                    *
     27 0B0000             START  ORG.L  $0B0000
     28 0B0000 A02D        START  DC     GETNXT
     29 0B0002 0C05 003F          CMP.B  #'?',D5
     30 0B0006 6700 00E4(B00EC)   BEQ    HELP
     31 0B000A 49FA 02EB(B02F4)   LEA    PROMPT(PC),A4
     32 0B000E A035               DC     PSTRNG         PROMPT FOR COMMAND
     33 0B0010 A029               DC     GETCH
     34 0B0012 A031               DC     TOUPPR         CONVERT TO UPPER CASE
```

```
35  080014 0C05 004E            CMP.B    #'N',D5
36  080018 6728       |B0042    BEQ.S    NEWPAG
37  08001A 0C05 0046            CMP.B    #'F',D5
38  08001E 6714       |B0034    BEQ.S    NEXTPG
39  080020 0C05 0042            CMP.B    #'B',D5
40  080024 6712       |B0038    BEQ.S    LASTPG
41  080026 0C05 0051            CMP.B    #'Q',D5
42  08002A 6706       |B0032    BEQ.S    EXIT
43  08002C 0C05 0045            CMP.B    #'E',D5
44  080030 66CE       |B0000    BNE.S    START
45  08032 A01E         EXIT     DC       WARMST
46                     *
47                     * COMMAND DISPATCHING
48                     *
49                     * NEXT PAGE
50  08034 6116     |B004C NEXTPG BSR.S   OPAGE        POINTER ALREADY SET FOR NEXT PAGE
51  08036 60C8     |B0000        BRA.S   START
52                     * LAST PAGE
53  08038 91FC 0000 0200 LASTPG SUB.L   #$0200,A0
54  0803E 610C     |B004C        BSR.S   OPAGE
55  08040 60BE     |B0000        BRA.S   START
56                     * NEW PAGE
57 >08042 6100 0064|B00A8 NEWPAG BSR     GETPG        NEED SUBR TO GET 4 HEX DIGITS
58  08046 2040                   MOVE.L  D0,A0
59  08048 6102     |B004C        BSR.S   OPAGE
60  0804A 60B4     |B0000        BRA.S   START
61                     *
62                     * ROUTINE TO OUTPUT A PAGE IN HEX AND ASCII
63                     *
64  0804C A034         OPAGE     DC      PCRLF
65  08004E 4240                  CLR.W   D0           LINE COUNTER
66  08050 2808                   MOVE.L  A0,D4
67  08052 A03C                   DC      OUTBH        PRINT PAGE ADDRESS
68  08054 A034                   DC      PCRLF
69                     * LOOP FOR LINES
70  08056 323C 000F    LLOOP     MOVE.W  #15,D1       COUNTER FOR CHARACTERS
71  0805A 1800                   MOVE.B  D0,D4
72  0805C E904                   ASL.B   #4,D4        ADDRESS OF FIRST BYTE OF LINE ON PAGE
73  0805E A03A                   DC      OUT2H
74  08060 183C 0020            MOVE.B  #$20,D4
75  08064 A033                   DC      PUTCH        SPACE
76  08066 A033                   DC      PUTCH        SECOND SPACE
77                     * INSIDE LOOP FOR 16 CHARACTERS IN HEX
78  08068 1818         CLOOP     MOVE.B  (A0)+,D4
79  0806A A03A                   DC      OUT2H        OUTPUT FIRST BYTE
80  0806C 1B3C 0020            MOVE.B  #$20,D4
81  08070 A033                   DC      PUTCH        SPACE
82  08072 57C9 FFF4|B0068       DBEQ    D1,CLOOP     CHARACTERS
83                     * NOW DO ASCII CHARACTERS, "." FOR NON PRINTABLE
84  08076 183C 0020            MOVE.B  #$20,D4
85  0807A A033                   DC      PUTCH        EXTRA SPACE BEFORE ASCII
86  0807C 91FC 0000 0010       SUB.L   #16,A0       BACK UP 16 LOCATIONS
87  08082 323C 000F            MOVE.W  #15,D1       RELOAD COUNTER FOR CHARACTERS
88  08086 1818         ALOOP     MOVE.B  (A0)+,D4
89  08088 0244 007F            AND     #$7F,D4      MASK OFF HI ORDER BIT
90  0808C 0C04 0020            CMP.B   #$20,D4      IS IT PRINTABLE?
91  08090 6C04     |B0096       BGE.S   AL1          IF YES
92  08092 183C 002E            MOVE.B  #'.',D4      ELSE PRINT PERIOD
93  08096 A033         AL1       DC      PUTCH
94  08098 57C9 FFEC|B0086       DBEQ    D1,ALOOP
95  0809C A034                   DC      PCRLF
96  0809E 5240                   ADD     #1,D0
97  0800A0 0C00 0010            CMP.B   #16,D0
98  0800A4 66B0     |B0056       BNE     LLOOP        NEXT LINE
99  0800A6 4E75                  RTS
100                    *
101                    * GET PAGE NUMBER IN A0 FOR N COMMAND
102                    *
103  0800A8 4281       GETPG     CLR.L   D1
104  0800AA 4280                 CLR.L   D0
105  0800AC A029       LOOP      DC      GETCH
106  0800AE A031                 DC      TOUPPR
107  0800B0 1005                 MOVE.B  D5,D0
108  0800B2 0C00 0030            CMP.B   #'0',D0      ASCII 0
109  0800B6 6D2C     |B00E4       BLT.S   ERR
110  0800B8 0C00 0039            CMP.B   #'9',D0      ASCII 9
111  0800BC 6E06     |B00C4       BGT.S   GETPG1
112  0800BE 0400 0030            SUB.B   #$30,D0      MAKE IT HEX
113  0800C2 6010     |B00D4       BRA.S   SHFT
114  0800C4 0C00 0041  GETPG1    CMP.B   #'A',D0      ASCII A
115  0800C8 6D1A     |B00E4       BLT.S   ERR
```

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PASC from S.E. Media - A FLEX9, SK*DOS Compiler with a definite Pascal "flavor". Anyone with a bit of Pascal experience should be able to begin using PASC to good effect in short order. The PASC package comes complete with three sample programs: ED (a syntax or structure editor), EDITOR (a simple, public domain, screen editor) and CHESS (a simple chess program). The PASC package comes complete with source (written in PASC) and documentation.
     *FLEX, SK*DOS $95.00*

WHIMSICAL from S.E. MEDIA Now supports *Real Numbers.* "Structured Programming" WITHOUT losing the Speed and Control of Assembly Language! Single-pass Compiler features unified, user-defined I/O; produces ROMable Code; Procedures and Modules (including pre-compiled Modules); many "Types" up to 32 bit Integers, 6-digit Real Numbers, unlimited sized Arrays (vectors only); Interrupt handling; long Variable Names; Variable Initialization; Include directive; Conditional compiling; direct Code insertion; control of the Stack Pointer, etc. Run-Time subroutines inserted as called during compilation. *Normally produces 10% less code than PL/9.*
     F, S and CCF - $195.00

KANSAS CITY BASIC from S.E. Media - *Basic for Color Computer OS-9* with many new commands and sub-functions added. A full implementation of the IF-THEN-ELSE logic is included, allowing nesting to 255 levels. Strings are supported and a subset of the usual string functions such as LEFT$, RIGHT$, MID$, STRING$, etc. are included. Variables are dynamically allocated. Also included are additional features such as Peek and Poke. A must for any Color Computer user running OS-9.
     CoCo OS-9 $39.95

C Compiler from Windrush Micro Systems by James McCosh. Full C for FLEX, SK*DOS except bit-fields, including an Assembler. *Requires the TSC Relocating Assembler if user desires to implement his own Libraries.*
     F, S and CCF - $295.00

C Compiler from Introl -- Full C except Doubles and Bit Fields, streamlined for the 6809. Reliable Compiler. FAST, efficient Code. More UNIX Compatible than most.
*FLEX, SK*DOS, CCF, OS-9 (Level II ONLY), U - $575.00*

PASCAL Compiler from Lucidata -- ISO Based P-Code Compiler. Designed especially for Microcomputer Systems. Allows linkage to Assembler Code for maximum flexibility.
     F, S and CCF 5" - $190.00   F, S 8"- $205.00

PASCAL Compiler from OmegaSoft (now Certified Software) -- For the *PROFESSIONAL;* ISO Based, Native Code Compiler. Primarily for Real-Time and Process Control applications. Powerful; Flexible.
     OS-9, F, S and CCF - $550.00
     OS-9 68000 Version - $900.00

KBASIC - from S.E. MEDIA -- A "Native Code" BASIC Compiler which is now Fully TSC XBASIC compatible. The compiler compiles to Assembly Language Source Code. A NEW, streamlined, Assembler is now included allowing the assembly of LARGE Compiled K-BASIC Programs. Conditional assembly reduces Run-time package.
*FLEX, SK*DOS, CCF, OS-9 Compiler /Assembler $99.00*

CRUNCH COBOL from S.E. MEDIA -- Supports large subset of ANSII Level I *COBOL* with many of the useful Level 2 features. Full FLEX, SK*DOS File Structures, including Random Files and the ability to process Keyed Files. Segment and link large programs at runtime, or implemented as a set of overlays. The System requires 56K and CAN be run with a single Disk System. *A very popular product.*
     *FLEX, SK*DOS, CCF - $99.95*

FORTH from Stearns Electronics -- A CoCo FORTH Programming Language. Tailored to the CoCo! Supplied on Tape, transferable to disk. Written in FAST ML. Many CoCo functions (Graphics, Sound, etc.). Includes an Editor, Trace, etc. Provides CPU Carry Flag accessibility, Fast Task Multiplexing, Clean Interrupt Handling, etc. for the "Pro". Excellent "Learning" tool!
     Color Computer ONLY - $58.95

FORTHBUILDER is a stand-alone target compiler (crosscompiler) for producing custom Forth systems and application programs. All of the 83-standard defining words and control structures are recognized by FORTHBUILDER.

FORTHBUILDER is designed to behave as much as possible like a resident Forth interpreter/compiler, so that most of the established techniques for writing Forth code can be used without change. Like compilers for other languages, FORTHBUILDER can operate in "batch mode".

The compiler recognizes and emulates target names defined by CONSTANT or VARIABLE and is readily extended with "compile-time" definitions to emulate specific target words.

FORTHBUILDER is supplied as an executable command file configured for a specific host system and target processor. Object code produced from the accompanying model source code is royalty-free to licensed users.
     F, CCF, S - $99.95

EDITORS & WORD PROCESSING

JUST from S.E. Media -- Text Formatter developed by Ron Anderson; for Dot Matrix Printers, provides many unique features. Output "Formatted" Text to the Display. Use the FPRINT.CMD supplied for producing multiple copies of the "Formatted" Text on the Printer INCLUDING IMBEDDED PRINTER COMMANDS (very useful at other times also, and worth the price of the program by itself). "User Configurable" for adapting to other Printers (comes set up for Epson MX-80 with Graftrax); up to ten (10) imbedded "Printer Control Commands". Compensates for a "Double Width" printed line. Includes the normal line width, margin, indent, paragraph, space, vertical skip lines, page length, page numbering, centering, fill, justification, etc. Use with PAT or any other editor.
* Now supplied as a two disk set:
Disk #1: JUST2.CMD object file.
*JUST2.TXT PL9 source: FLEX, SK*DOS - CC*
Disk #2: JUSTSC object and source in C:
*FLEX, SK*DOS - OS9 - CC*
The JTSC and regular JUST C source are two separate programs. JTSC compiles to a version that expects TSC Word Processor type commands, (.pp .sp .ce etc.) Great for your older text files. The C

ITS EASY TO USE!

XDMS-IV keeps data management simple! Rather than design a complex DBMS which hides the true nature of the data, we kept XDMS-IV file oriented. The user view of data relationships is presented in reports and screen output, while the actual data resides in easy to maintain files. This aspect permits customized presentation and reports without complex redefinition of the database files and structure. XDMS-IV may be used for a wide range of applications from simple record management systems (addresses, inventory ...) to integrated database systems (order entry, accounting...)

The possibilities are unlimited...

FOR 6809 FLEX-SK*DOS(5/8")   **$249.95**

UTILITIES

Basic09 XRef from S.E. Media -- This Basic09 Cross Reference Utility is a Basic09 Program which will produce a "pretty printed" listing with each line numbered, followed by a complete cross referenced listing of all variables, external procedures, and line numbers called. Also includes a Program List Utility which outputs a fast "pretty printed" listing with line numbers. Requires Basic09 or RunB.

  O & CCO obj. only -- $39.95; w/ Source - $79.95

BTree Routines - Complete set of routines to allow simple implementation of keyed files - *for your programs* - running under Basic09. A real time saver and should be a part of every serious programmers tool-box.

  O & CCO obj. only - $89.95

Lucidata PASCAL UTILITIES (Requires Pascal ver 3)

XREF -- produce a Cross Reference Listing of any text; oriented to Pascal Source.

INCLUDE -- Include other Files in a Source Text, including Binary - unlimited nesting.

PROFILER -- provides an Indented, Numbered, "Structogram" of a Pascal Source Text File; view the overall structure of large programs, program integrity, etc. Supplied in Pascal Source Code; requires compilation.

  F, S, CCF --- EACH   5" - $40.00, 8" - $50.00

DUB from S.E. Media -- A UniFLEX BASIC **decompiler** Re-Create a Source Listing from UniFLEX Compiled basic Programs. Works w/ ALL Versions of 6809 UniFLEX basic.

  U - $219.95

LOW COST PROGRAM KITS from Southeast Media The following kits are available for FLEX, SK*DOS on either 5" or 8" Disk.

1.   **BASIC TOOL-CHEST $29.95**
   BLISTER.CMD: pretty printer
   LINEXREF.BAS: line cross-referencer
   REMPAC.BAS, SPCPAC.BAS, COMPAC.BAS:
   remove superfluous code
   STRIP.BAS: superfluous line-numbers stripper

2.   **FLEX, SK*DOS UTILITIES KIT $39.99**
   CATS. CMD: alphabetically-sorted directory listing
   CATD.CMD: date-sorted directory listing
   COPYSORT.CMD: file copy, alphabetically
   COPYDATE.CMD: file copy, by date-order
   FILEDATE.CMD: change file creation date
   INFO.CMD (& INFOGMX.CMD): tells disk attributes & contents
   RELINK.CMD (& RELINK82): re-orders fragmented free chain
   RESQ.CMD: undeletes (recovers) a deleted file
   SECTORS.CMD: show sector order in free chain
   XL.CMD: super text lister

3.   **ASSEMBLERS/DISASSEMBLERS UTILITIES $39.95**
   LINEFEED.CMD: 'modularise' disassembler output
   MATH.CMD: decimal, hex, binary, octal conversions
   & tables
   SKIP.CMD: column stripper

4.   **WORD - PROCESSOR SUPPORT UTILITIES $49.95**
   FULLSTOP.CMD: checks for capitalization
   BSTYCIT.BAS (.BAC): Stylo to dot-matrix printer
   NECPRINT.CMD: Stylo to dot-matrix printer filter code

5.   **UTILITIES FOR INDEXING $49.95**
   MENU.BAS: selects required program from list below
   INDEX.BAC: word index
   PHRASES.BAC: phrase index
   CONTENT.BAC: table of contents
   INDXSORT.BAC: fast alphabetic sort routine
   FORMATER.BAC: produces a 2-column formatted index
   APPEND.BAC: append any number of files
   CHAR.BIN: line reader

BASIC09 TOOLS consist of 21 subroutines for Basic09.
6 were written in C Language and the remainder in assembly.
All the routines are compiled down to native machine code which makes them fast and compact.

   1. **CFILL** -- fills a string with characters
   2. **DPEEK** -- Double peek
   3. **DPOKE** -- Double poke
   4. **FPOS** -- Current file position
   5. **FSIZE** -- File size
   6. **FTRIM** -- removes leading spaces from a string
   7. **GETPR** -- returns the current process ID
   8. **GETOPT** -- gets 32 byte option section
   9. **GETUSR** -- gets the user ID
   10. **GTIME** -- gets the time
   11. **INSERT** -- insert a string into another
   12. **LOWER** -- converts a string into lowercase
   13. **READY** -- Checks for available input
   14. **SETPRIOR** -- changes a process priority
   15. **SETUSR** -- changes the user ID
   16. **SETOPT** -- set 32 byte option packet
   17. **STIME** -- sets the time
   18. **SPACE** -- adds spaces to a string
   19. **SWAP** -- swaps any two variables
   20. **SYSCALL** -- system call
   21. **UPPER** -- converts a string to uppercase

For OS-9 - $44.95 - Includes Source Code
   Limited Special - $19.95

SOFTOOLS

The following programs are included in object form for immediate application. PL/9 source code available for customization.

READ-ME Complete instructions for initial set-up and operation. Can even be printed out with the included text processor.

CONFIG one time system configuration.

CHANGE changes words, characters, etc. globally to any text type file.

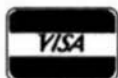**CLEANTXT** converts text files to standard FLEX, SK*DOS files.

COMMON compare two text files and reports differences.

COMPARE another check file that reports mis-matched lines.

CONCAT similar to FLEX, SK*DOS append but can also list files to screen.

DOCUMENT for PL/9 source files. Very useful in examining parameter passing aspects of procedures.

COPYCAT from Lucidata -- *Pascal NOT required.* Allows reading TSC Mini-FLEX, SK*DOS, SSB DOS68, and Digital Research CP/M Disks while operating under SK*DOS , FLEX1.0, FLEX 2.0, or FLEX 9.0 with 6800 or 6809 Systems. COPYCAT will not perform miracles, but, between the program and the manual, you stand a good chance of accomplishing a transfer. Also includes some Utilities to help out. Programs supplied in Modular Source Code (Assembly Language) to help solve unusual problems.

   F, S and CCF 5" - $50.00   F, S 8" - $65.00

VIRTUAL TERMINAL from S.E. Media - Allows one terminal to do the work of several. The user may start as many as eight tasks on one terminal, under *VIRTUAL TERMINAL* and switch back and forth between tasks at will. No need to exit each one; just jump back and forth. Complete with configuration program. The best way to keep up with those background programs.

   6809 O & CCO - obj. only - $49.95

FLEX, SK*DOS DISK UTILITIES from Computer Systems Consultants -- Eight (8) different Assembly Language (w/ Source Code) FLEX, SK*DOS Utilities for every FLEX, SK*DOS Users Toolbox: Copy a File with CRC Errors; Test Disk for errors; Compare two Disks; a fast Disk Backup Program; Edit Disk Sectors; Linearize Free-Chain on the Disk; print Disk Identification; and Sort and Replace the Disk Directory (in sorted order). -- PLUS -- Ten XBASIC Programs including: A BASIC Resequencer with EXTRAs over "RENUM" like check for missing label definitions, processes Disk to Disk instead of in Memory, etc. Other programs Compare, Merge, or Generate Updates between two BASIC Programs, check BASIC Sequence Numbers, compare two unsequenced files, and 5 Programs for establishing a Master Directory of several Disks, and sorting, selecting, updating, and printing paginated listings of these files. A BASIC Cross-Reference Program, written in Assembly Language, which provides an X-Ref Listing of the Variables and Reserved Words in TSC BASIC, XBASIC, and PRECOMPILER BASIC Programs.

ALL Utilities include Source (either BASIC or A.L. Source Code).
   F, S and CCF - $50.00
   BASIC Utilities ONLY for UniFLEX -- $30.00

MS-DOS-FLEX Transfer Utilities to OS-9 For 68XXX and CoCo* OS-9 Systems Now READ - WRITE - DIR - DUMP - EXPLORE FLEX & MS-DOS Disk. These Utilities come with a rich set of options allowing the transfer of text type files from/to FLEX & MS-DOS disks. *CoCo systems require the D.P. Johnson SDISK utilities and OS-9 and two drives of which one must be a "host" floppy.

   *CoCo Version: $69.95   68XXX Version $99.95

MISCELLANEOUS

TABULA RASA SPREADSHEET from Computer Systems Consultants -- TABULA RASA is similar to DESKTOP/PLAN; provides use of tabular computation schemes used for analysis of business, sales, and economic conditions. Menu-driven; extensive report-generation capabilities. Requires TSC's Extended BASIC.
   F, S and CCF , U - $50.00, w/ Source - $100.00

DYNACALC -- Electronic Spread Sheet for the 6809 and 68000.
   F, S, OS-9 and SPECIAL CCF - $200.00,   U - $395.00
   OS-9 68K - $595.00

FULL SCREEN INVENTORY/MRP from Computer Systems Consultants Use the Full Screen Inventory System/Materials Requirement Planning for maintaining inventories. Keeps item field file in alphabetical order for easier inquiry. Locate and/or print records matching partial or complete item, description, vendor, or attributes; find backorder or below stock levels. Print-outs in item or vendor order. MRP capability for the maintenance and analysis of Hierarchical assemblies of items in the inventory file. Requires TSC's Extended BASIC.
   F, S and CCF, U - $50.00, w/ Source - $100.00

FULL SCREEN MAILING LIST from Computer Systems Consultants -- The Full Screen Mailing List System provides a means of maintaining simple mailing lists. Locate all records matching on partial or complete name, city, state, zip, or attributes for Listings or Labels, etc. Requires TSC's Extended BASIC.
   F, S and CCF, U - $50.00, w/ Source - $100.00

DIET-TRAC Forecaster from S.E. Media -- An XBASIC program that plans a diet in terms of either calories and percentage of carbohydrates, proteins and fats (C P G%) or grams of Carbohydrate, Protein and Fat food exchanges of each of the six basic food groups (vegetable, bread, meat, skim milk, fruit and fat) for a specific individual. Sex, Age, Height, Present Weight, Frame Size, Activity Level and Basal Metabolic Rate for nonnal individual are taken into account. Ideal weight and sustaining calories for any weight of the above individual are calculated. Provides number of days and daily calendar after weight goal and calorie plan is determined.
   F, S - $59.95,   U - $89.95

GAMES

RAPIER - 6809 Chess Program from S.E. Media -- Requires FLEX, SK*DOS and Displays on Any Type Terminal. Features: Four levels of play. Swap side. Point scoring system. Two display boards. Change skill level. Solve Checkmate problems in 1-2-3-4 moves. Make move and swap sides. Play white or black. This is one of the strongest CHESS programs running on any microcomputer, *estimated USCF Rating 1600+ (better than most 'club' players at higher levels)*
   F, S and CCF - $79.95

NEW

MS-DOS/FLEX Transfer Utilities For 68XXX and CoCo* OS-9 Systems. Now Read, Write, DIR, Dump and Explore FLEX & MS-DOS Disks. Supplied with a rich set of options to explore and transfer text type files from/to FLEX and MS-DOS disks. *CoCo OS-9 requires SDISK utilities & two floppy drives.
   CCO $69.95   68XXX OS-9 $99.95

```
116  0B00CA 0C00 0046            CMP.B    #'F',D0        ASCII F
117  0B00CE 6E14    {B00E4       BGT.S    ERR
118  0B00D0 0400 0037            SUB.B    #$37,D0
119  0B00D4 E900        SHFT     ASL.B    #4,D0          000000B0
120  0B00D6 E980                 ASL.L    #4,D0          00000B00
121  0B00D8 5241                 ADD      #1,D1
122  0B00DA 0C01 0004            CMP.B    #4,D1
123  0B00DE 66CC    {800AC       BNE.S    LOOP
124  0B00E0 A034                 DC       PCRLF
125  0B00E2 4E75                 RTS
126                          *
127  0B00E4 49FA 01FA{B02E0 ERR  LEA      MESG(PC),A4
128  0B00E8 A035                 DC       PSTRNG
129  0B00EA 4E75                 RTS
130                          * HELP MESSAGE PRINTS OUT ON COMMAND: MDUMP ?
131                          *
132  0B00EC 49FA 0006{B00F4 HELP LEA      HLPMSG(PC),A4
133  0B00F0 A035                 DC       PSTRNG
134  0B00F2 A01E                 DC       WARMST
135                          *
136  0B00F4 4D44 554D 5020 HLPMSG DC.B    "MDUMP requires no command line arguments.  It loads",$0D,$0A
137  0B0129 696E 2068 6967        DC.B    "in high memory, presently $0B0000 plus OFFSET.  It prompts",$0D,$0A
138  0B0165 666F 7220 6120        DC.B    "for a command.  N005E <cr> will cause it to dump a page ",$0D,$0A
139  0B019F 6F66 2060 656D        DC.B    "of memory, in this case from $005E00 to $005EFF.",$0D,$0A
140  0B01D1 636F 6D6D 616E        DC.B    "command F will cause the next page (forward) to be dumped",$0D,$0A
141  0B020C 616E 6420 4220        DC.B    "and B will cause the previous page (backward) to be dumped",$0D,$0A
142  0B0248 4561 6368 2070        DC.B    "Each page is dumped 16 bytes to a line with the ASCII",$0D,$0A
143  0B027F 7265 7072 6573        DC.B    "representation following the HEX.  Non-printable ASCII",$0D,$0A
144  0B02B7 6368 6172 6163        DC.B    "characters are represented as periods.",$0D,$0A,$04
145
146  0B02E0 4E4F 5420 4120 MESG   DC.B    "NOT A VALID ADDRESS",$04
147  0B02F4 434F 4D4D 414E PROMPT DC.B    "COMMAND: ",$04
148                               END     START
   0 ERRORS DETECTED


   1
   2
   3
   4              ******************************************************
   5              * DISKFILE DUMP PROGRAM                              *
   6              *                                                    *
   7              * COMMANDS:                                          *
   8              * N XXXX NEXT PAGE TO BE DUMPED                      *
   9              * B BACK A SECTOR                                    *
  10              * F FORWARD A SECTOR                                 *
  11              ******************************************************
  12
  13              * SK*DOS / 68K EQUATES FOR USER PROGRAMS
  14
  15       0000A029    GETCH   EQU    $A029        Get input character with echo (7 bits)
  16       0000A023    GETNAM  EQU    $A023        Get filename into FCB
  17       0000A024    DEFEXT  EQU    $A024        Set default extension
  18       0000A005    FOPENR  EQU    $A005        Open file for read
  19       0000A001    FREAD   EQU    $A001        Read a byte
  20       0000A011    FSKIP   EQU    $A011        Next Sector
  21       0000A031    TOUPPR  EQU    $A031        Conv't char in D5 to Upper Case
  22       0000A02F    HEXIN   EQU    $A02F        Input hexadecimal number
  23       0000A03A    OUT2H   EQU    $A03A        Output 2 hex digits
  24       0000A03B    OUT4H   EQU    $A03B
  25       0000A03C    OUT8H   EQU    $A03C        Output 8 hex digits
  26       0000A034    PCRLF   EQU    $A034        Print CR/LF
  27       0000A036    PNSTRN  EQU    $A036        Print string (without CR/LF)
  28       0000A035    PSTRNG  EQU    $A035        Print CR/LF and string
  29       0000A033    PUTCH   EQU    $A033        Output character
  30       0000A000    VPOINT  EQU    $A000        Point to SK*DOS variable area
  31       0000A01E    WARMST  EQU    $A01E        Warm start
  32
  33                       *
  34                       *
  35  000000 A000    START  DC     VPOINT        GET POINTER
  36  000002 284E           MOVE.L A6,A4          FCB POINTER
  37  000004 A023           DC     GETNAM        GET FILENAME
  38  000006 65D0 0096{0009E BCS    HELP
  39  00000A 7801           MOVE.L #1,D4          CODE FOR .TXT
  40  00000C A024           DC     DEFEXT
  41  00000E A005           DC     FOPENR
  42  000010 43FA 0032{00044 LEA   OLDTS(PC),A1
  43  000014 328C 0003      MOVE.W #$0003,(A1)    PRIME WITH BAD TS
  44  000018 A001    LOOP   DC     FREAD
  45  00001A 6526    {00042  BCS.S ERROR
```

```
46  00001C 204C                      MOVE.L   A4,A0
47  00001E D1FC 0000 0060            ADD.L    #96,A0          POINT AT SECTOR INFO
48  000024 43FA 001E(00044           LEA      OLDTS(PC),A1
49  000028 3011                      MOVE.W   (A1),D0
50  00002A B050                      CMP.W    (A0),D0
51  00002C 6714     (00042           BEQ.S    ERROR           END OF FILE
52  00002E 3290                      MOVE.W   (A0),(A1)
53  000030 6114     (00046           BSR.S    OPAGE
54  000032 A029                      DC       GETCH
55  000034 A031                      DC       TOUPPR
56  000036 0C05 0045                 CMP.B    #'E',05
57  00003A 6602     (0003E           BNE.S    CONTIN
58  00003C A01E                      DC       WARMST
59  00003E A011            CONTIN    DC       FSKIP
60  000040 6006     (00018           BRA      LOOP
61  000042 A01E            ERROR     DC       WARMST          ASSUME END OF FILE
62  000044                 OLDTS     OS.W     1
63                      *
64                      * ROUTINE TO OUTPUT A PAGE IN HEX AND ASCII
65                      *
66  000046 A034            OPAGE     DC       PCRLF
67  000048 4240                      CLR.W    D0              LINE COUNTER
68  00004A A034                      DC       PCRLF
69                      * LOOP FOR LINES
70  00004C 323C 000F       LLOOP     MOVE.W   #15,D1          COUNTER FOR CHARACTERS
71  000050 1800                      MOVE.B   D0,D4
72  000052 E904                      ASL.B    #4,D4           ADDRESS OF FIRST BYTE OF LINE ON PAGE
73  000054 A03A                      DC       OUT2H
74  000056 183C 0020                 MOVE.B   #$20,D4
75  00005A A033                      DC       PUTCH           SPACE
76  00005C A033                      DC       PUTCH           SECOND SPACE
77                      * INSIDE LOOP FOR 16 CHARACTERS IN HEX
78  00005E 1818            CLOOP     MOVE.B   (A0)+,D4
79  000060 A03A                      DC       OUT2H           OUTPUT FIRST BYTE
80  000062 183C 0020                 MOVE.B   #$20,D4
81  000066 A033                      DC       PUTCH           SPACE
82  000068 57C9 FFF4(0005E            DBEQ     D1,CLOOP        CHARACTERS
83                      * NOW DO ASCII CHARACTERS, "." FOR NON PRINTABLE
84  00006C 183C 0020                 MOVE.B   #$20,D4
85  000070 A033                      DC       PUTCH           EXTRA SPACE BEFORE ASCII
86  000072 91FC 0000 0010            SUB.L    #16,A0
87  000078 323C 000F                 MOVE.W   #15,D1          RELOAD COUNTER FOR CHARACTERS
88  00007C 1818            ALOOP     MOVE.B   (A0)+,D4
89  00007E 0244 007F                 AND      #$7F,D4         MASK OFF HI ORDER BIT
90  000082 0C04 0020                 CMP.B    #$20,D4         IS IT PRINTABLE?
91  000086 6C04     (0008C           BGE.S    AL1             IF YES
92  000088 183C 002E                 MOVE.B   #'.',D4         ELSE PRINT PERIOD
93  00008C A033            AL1       DC       PUTCH
94  00008E 57C9 FFEC(0007C            DBEQ     D1,ALOOP
95  000092 A034                      DC       PCRLF
96  000094 5240                      ADD      #1,D0
97  000096 0C00 0010                 CMP.B    #16,D0
98  00009A 66B0     (0004C           BNE      LLOOP           NEXT LINE
99  00009C 4E75                      RTS
100
101 00009E 49FA 0006(000A6 HELP      LEA      HLPMSG(PC),A4
102 0000A2 A035                      DC       PSTRNG
103 0000A4 A01E                      DC       WARMST
104
105 0000A6 5379 6E74 6178  HLPMSG    DC.B     "Syntax: DDUMP FILENAME   Defaults are work drive",$0D,$0A
106 0000D8 616E 6420 2E54            DC.B     "and .TXT extension.",$00,$0A,$0A
107 0000EE 4444 554D 5020            DC.B     "DDUMP dumps a disk file to the terminal one sector",$00,$0A
108 000122 6174 2061 2074            DC.B     "at a time.  The dump displays 16 bytes in HEX followed by",$0D,$0A
109 00015D 7468 6520 4153            DC.B     "the ASCII representation of the same 16 bytes.  Non-",$0D,$0A
110 000193 7072 696E 7461            DC.B     "printable characters are displayed as periods.",$0D,$0A
111 0001C3 4174 2074 6865            DC.B     "At the command prompt, F will cause the next (forward)",$0D,$0A
112 0001FB 7365 6374 6F72            DC.B     "sector to be displayed, and B (back) will display the",$0D,$0A
113 000232 7072 6576 696F            DC.B     "previous sector.  E will Exit the program as will",$0D,$0A
114 000265 4620 7768 656E            DC.B     "F when the last sector of the file has been reached.",$0D,$0A,$04
115
116                        END       START

  0 ERRORS DETECTED
```

```
   1                         * FIND UTILITY FOR SK*DOS /68K
   2                         * BY R. W. ANDERSON 1988
   3                         *
   4                         * EQUATES TO SK*DOS
   5                         *
   6                         * THE FOLLOWING WOULD NORMALLY BE IN SKEQUATE.TXT
   7          0000A000       VPOINT   EQU     $A000
   8          0000A038       OUT5D    EQU     $A038
   9          00000001       FCBERR   EQU     1
  10          0000A024       DEFEXT   EQU     $A024
  11          0000A008       FCLOSE   EQU     $A008
  12          0000A005       FOPENR   EQU     $A005
  13          0000A006       FOPENW   EQU     $A006
  14          0000A001       FREAD    EQU     $A001
  15          0000A002       FWRITE   EQU     $A002
  16          0000A023       GETNAM   EQU     $A023
  17          0000A029       GETCH    EQU     $A029
  18          0000A035       PSTRNG   EQU     $A035
  19          0000A036       PNSTRN   EQU     $A036
  20          0000A034       PCRLF    EQU     $A034
  21          0000A037       PERROR   EQU     $A037
  22          0000A033       PUTCH    EOU     $A033
  23          0000A01E       WARMST   EQU     $A01E
  24                         *
  25  000000                          ORG     $0000
  26  000000 6002    {00004 FIND       BRA.S   START               GOTO START
  27                         *
  28  000002 0100             VER      DC.W    $0100               VERSION NUMBER
  29                         *
  30  000004 A034             START    DC      PCRLF               START ON NEW LINE
  31  000006 204E                      MOVE.L  A6,A0               SAVE POINTER
  32  000008 284E                      MOVE.L  A6,A4               POINTER TO USER FCB
  33  00000A A023                      DC      GETNAM              GET FILE SPEC
  34  00000C 6404    {00012            BCC.S   NAMEOK              IF FILENAME OK
  35  00000E 6000 0006{000E6           BRA     HELP
  36                         * IF FILE SPEC WAS OK; DEFAULT TO .TXT
  37  000012 183C 0001        NAMEOK   MOVE.B  #1,04               DEFAULT EXTENSION
  38  000016 A024                      DC      DEFEXT              DEFAULT EXTENSION
  39                         *
  40                         * NOW OPEN THE FILE.
  41  000018 A005                      DC      FOPENR              OPEN FOR READ
  42  00001A 6660    {0007C            BNE.S   ERROR1              IF NOT ZERO
  43                         *
  44                         * MAIN LOOP TO READ LINES AND SEARCH FOR MATCH
  45                         *
  46  00001C 49FA 01DE{001FC MAIN       LEA     PROMPT(PC),A4
  47  000020 A035                       DC      PSTRNG              -INPUT SEARCH STRING-
  48  000022 41FA 02EE{00312            LEA     LINECT(PC),A0
  49  000026 30BC 0000                  MOVE.W  #0,(A0)             CLEAR LINE COUNTER
  50  00002A 6158    {00084             BSR.S   GETSTR              GET SEARCH STRING INTO BUFFER
  51  00002C A034                       DC      PCRLF
  52                         *
  53  00002E A000             LINE      DC      VPOINT
  54  000030 43FA 01E0{00212           LEA     LINBUF(PC),A1
  55  000034 284E             LINE1     MOVE.L  A6,A4               POINT TO SYSFCB
  56  000036 A001                       DC      FREAD               GO READ NEXT CHAR
  57  000038 6636    {00070             BNE.S   ERROR
  58                         *
  59  00003A 12C5                        MOVE.B  D5,(A1)+            PUT IN LINE BUFFER
  60  00003C 0C05 0000                  CMP.B   #$0D,D5             IS IT LINEFEED
  61  000040 66F2    {00034             BNE.S   LINE1               GET MORE
  62  000042 137C 0004 FFFF             MOVE.B  #$04,-1(A1)         CHANGE CR TO $04
  63  000048 41FA 02C8{00312            LEA     LINECT(PC),A0
  64  00004C 5250                        ADD.W   #1,(A0)
  65  00004E 6160    {000B0             BSR.S   MATCH
  66  000052 4A00                        TST.B   D0
  67  000052 6602    {00056             BNE.S   FOUND
  68  000054 60D8    {0002E             BRA.S   LINE
  69  000056 383A 028A{00312 FOUND      MOVE.W  LINECT(PC),D4
  70  00005A 1A3C 00FF                   MOVE.B  #-1,D5              SET NON-ZERO
  71  00005E A038                        DC      OUT5D
  72  000060 183C 0020                   MOVE.B  #$20,D4
  73  000064 A033                        DC      PUTCH
  74  000066 49FA 01AA{00212            LEA     LINBUF(PC),A4
  75  00006A A036                        DC      PNSTRN
  76  00006C A034                        DC      PCRLF
  77  00006E 60BE    {0002E             BRA.S   LINE
  78                         *
  79  000070 0C2C 0008 0001 ERROR .      CMP.B   #8,FCBERR(A4)
  80  000076 6604    {0007C             BNE.S   ERROR1              NOT END OF FILE
  81  000078 6106    {000B0 ERCLS       BSR.S   CLOSE
  82  00007A A01E                        DC      WARMST
```

```
 83   00007C A037              ERROR1  DC       PERROR
 84   00007E 60F8     100078           BRA.S    ERCLS
 85                         *
 86                         * SUBROUTINES
 87                         *
 88   000080 A008              CLOSE   DC       FCLOSE           CLOSE FILE
 89   000082 4E75              RTS
 90                         *
 91                         * SUBROUTINE TO GET A STRING FROM TERMINAL INTO
 92                         * A BUFFER USING PC RELATIVE ADDRESSING
 93                         *
 94   000084 41FA 020C|00292 GETSTR  LEA      BUFF(PC),A0      GET POINTER TO BUFFER
 95   000088 A029              GET1    DC       GETCH
 96   00008A 0C05 000D                 CMP.B    #$0D,D5
 97   00008E 671A     |000AA           BEQ.S    EXIT
 98   000090 0C05 0008                 CMP.B    #$08,D5
 99   000094 6610     |000A6           BNE.S    GET2
100   000096 5388                      SUB.L    #1,A0
101   000098 183C 0020                 MOVE.B   #$20,D4
102   00009C A033                      DC       PUTCH
103   00009E 183C 0008                 MOVE.B   #$08,D4
104   0000A2 A033                      DC       PUTCH
105   0000A4 60E2     |00088           BRA.S    GET1
106   0000A6 10C5              GET2    MOVE.B   D5,(A0)+
107   0000A8 60DE     |00088           BRA.S    GET1
108   0000AA 108C 0000        EXIT    MOVE.B   #0,(A0)          NULL TERMINATE STRING
109   0000AE 4E75              RTS
110                         *
111                         * SUBROUTINE TO SEEK A MATCH BETWEEN A SEARCH STRING
112                         * AND A LINE OF TEXT.  RETURNS ZERO IN D0 IF NOT FOUND,
113                         * 1 IF FOUND.
114                         *
115   0000B0 41FA 01E0|00292 MATCH   LEA      BUFF(PC),A0
116   0000B4 43FA 015C|00212         LEA      LINBUF(PC),A1
117   0000B8 2449                     MOVE.L   A1,A2           A2 KEEPS TRACK OF START OF MATCH
118   0000BA 1211              MATCH1  MOVE.B   (A1),D1         GET CHAR IN D1
119   0000BC B210                     CMP.B    (A0),D1
120   0000BE 6606     |000C6         BNE.S    MATCH2          IF NOT EQUAL MOVE DOWN LINE
121   0000C0 5288                     ADD.L    #1,A0           IF EQUAL COMPARE NEXT CHAR OF BUFF
122   0000C2 5289                     ADD.L    #1,A1           NEXT CHAR OF LINE
123   0000C4 6010     |000D6         BRA.S    MATCH3
124   0000C6 528A              MATCH2  ADD.L    #1,A2           IF NOT EQUAL START AGAIN
125   0000C8 224A                     MOVE.L   A2,A1           START AT NEXT CHAR IN LINE
126   0000CA 0C11 0004                 CMP.B    #$04,(A1)
127   0000CE 6712     |000E2         BEQ.S    NOTFND          GOT TO END OF LINE
128   0000D0 41FA 01C0|00292         LEA      BUFF(PC),A0
129   0000D4 60E4     |000BA         BRA.S    MATCH1          GO AROUND AGAIN.
130   0000D6 0C10 0000        MATCH3  CMP.B    #0,(A0)          MATCH
131   0000DA 66DE     |000BA         BNE.S    MATCH1
132   0000DC 103C 0001                 MOVE.B   #1,D0
133   0000E0 4E75              RTS
134   0000E2 4200              NOTFND  CLR.B    D0
135   0000E4 4E75              RTS
136                         *
137   0000E6 49FA 0006|000EE HELP    LEA      HLPMSG(PC),A4
138   0000EA A035              DC       PSTRNG
139   0000EC A01E              DC       WARMST
140                         *
141   0000EE 5379 6E74 6178   HLPMSG  DC.B     "Syntax:  FIND FILENAME    Defaults are work drive",$0D,$0A
142   000121 2020 2020 2020           DC.B     "                 and .TXT extension.",$0D,$0A
143   000142 4649 4E44 2070           DC.B     "FIND prompts for a search string.  It will search the file",$0D,$0A
144   00017E 616E 6420 6C69           DC.B     "and list each line in the file that contains the string.",$0D,$0A
145   0001B8 5468 6520 6C69           DC.B     "The line number is included for convenience in editing",$0D,$0A
146   0001F0 7468 6520 6669           DC.B     "the file.",$0D,$0A,$04
147
148   0001FC 494E 5055 5420   PROMPT  DC.B     "INPUT SEARCH STRING: ",$04
149   000212                  LINBUF  DS.B     128
150   000292                  BUFF    DS.B     128
151   000312                  LINECT  DS.W     1
152                            END      FIND
      0 ERRORS DETECTED


EOF
```

FOR THOSE WHO NEED TO KNOW

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A Review of Curator

An Art Management and Integration Program from Solutions International, Inc.

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What Is The Problem to be Solved?

In the beginning, there was MacPaint and McWrite and the transfer of images between programs was simple. We understood that MacPaint produced bitmapped images in the PNTS format which when cut or copied, assumed the clipboard's PICT format for transfer. McWrite saved its documents in the WORD format or as plain ASCI code. Transfers of images took place effortlessly. That was all we needed to know.

The world has changed! The Macintosh environment may be the simplest now available, but even so, the explosion of second and third generation software has resulted in a diverse population of image formats, few of which are compatible with all applications. Postscript output drives postscript printers like the Apple LaserWriter. Encapsulated postscript contains the postscript code and a Quick-Draw kernel that enables an approximation of the postscript image to be seen on the screen. The TIFF format is a high resolution bitmapped image commonly used with scanners. While TIFF is supposedly a standard format, I heard an industry insider claim that he personally knew of 47 versions of this format. A new "standard" format has just been introduced called RIFF. This format is reportedly much like TIFF, but requires only 30% of the storage space for an equivalent image. In addition to the large number of so called standard formats, there are a multitude of proprietary formats.

If we worked exclusively in one application, this would not be much of a problem. The real world is, however, that we use multiple applications using multiple output formats and with the obvious need to transfer images between programs. Programmers have taken some steps toward minimizing this problem. For example, SuperPaint documents can be saved in a PNTG (MacPaint) or PICT (Clipboard) format in addition to the proprietary format normally used by SuperPaint. As a result, SuperPaint images can be can figured for transfer into applications that only accept the PNTG format.

Such steps have not entirely solved the problem. For example, MacWrite does not understand Encapsulated PostScript and McDraw cannot read TIFF.

Then Along Came Curator

Curator allows you to have immediate access to the graphics you need regardless of what folder it is in. You do not have to set up a special art library. Curator can handle a variety of formats including PICT, MacPaint, TIFF, Postscript, Mac-Encapsulated Postscript, IBM-Encapsulated Postscript, plain postscript, and Glue. Then once you have found it, Curator converts the graphics from the format they were created in to the format you need.

Curator is provided in both application and desk accessory form. Both work essentially the same. Also provided is a small application called Curator's Assistant. This application automatically scans an entire floppy or hard disk and creates the Curator catalog with thumbnail sketches and key words. It is used when you first buy Curator to set up the initial catalog or to update the catalog later if you acquired a large quantity of additional graphics.

Getting Started

Drag the Curator and Curator Assistant Icons to your disks and load the CuratorDA and you are ready to go. (Note that you may, but you do not have to load both the application and the DA version of Curator since both do the same thing.) Double click the Curator Assistant Icon, respond to a dialog box, then work begins on constructing the Curator catalog. The manual warns you that this can take a long time. The catalog for my approximately 4 megabytes of graphics took less than 10 minutes.

Using Curator to Manage Art Collections

Curator allows you to search for and select graphics by name, thumbnail sketch, or key word.

Choosing the **Select by Name** option presents a typical dialog box used to select and open applications or documents. You can step through the various layers of folders then documents until you find what you want. You may select a different drive and eject a disk from this window. When you single click a document listed in the window, a 1" by 1-1/4" thumbnail sketch is displayed, the graphics format is indicated as is the application which created it. (Warning—These thumbnail sketches are so small that they do not effectively show some images. If you have a 8-1/2" by 11" image containing a number of small objects, as often occurs with commercial clip art, the thumbnail may not be recognizable. Of course, for bigger objects, the thumbnail sketches are helpful.) In any of your searches for graphics, you may limit the search to only a particular format or combination of formats. For example, you may limit your search to MacPaint and TIFF images thus excluding Mac EPS, IBM EPS, PICT, Postscript, and Glue.

The **Select by Thumbnail** option works essentially like the **Select by Name** option. Instead of a list of document names, you are presented with a window showing up to 10 thumbnails. You may scroll to see other thumnails if more than 10 documents are in the folder. Double clicking a thumbnailcauses that image to be displayed full size. Both

the **Select by Name** and the **Select by Thumbnail** look at only one folder at a time. They cannot be used to do a global search of a volume.

The **Search by Name** option is different from the options listed above in that it searches the entire volume. (I think the choice of terminology in this program is unduly confusing. It is not intuitive why one option is called **Search by Name** and the other **Select by Name**.) You enter all or part of a document named and any matches will be listed. Also, you may search by format. For example, you could look for all TIFF graphics regardless of its name. A bar graph keeps you informed of the status (i.e., percent complete) of the search.

There seems to be a flaw in the design of this feature. A search will often yield multiple 'hits'. Suppose that you want to look at a full-sized image of each hit to determine which one you want. After you view the first, you close the image window and return to the list of hits for further exploration, right? Wrong! Your choices have been erased and you have to start over. This does not make sense.

Curator allows you to assign one or more keywords to each image on your floppy or hard disk. It also allows you to use the keywords which are sometimes preassigned by the publishers of clip art or the key words in PictureBase files. The lists of keywords can be viewed and changed at any time. This feature can provide tremendous power in finding exactly the desired image quickly. It only works, however, if you are disciplined enough to assign meaningful keywords for each image.

Curator creates an index of all keywords on the entire floppy or hard disk. When you select **Search by Keyword**, a scrolling list of available key words is presented. You may then select on or more of these for inclusion in your search criteria.

The final search option provided by Curator is called **Browae**. This option allows you to step through the images of a selected folder to view a small portion of the image, assign or change keywords, or select and open documents. The image is viewed at full size through a 5" by 2" window. Moving form one image to another takes 5 or 6 seconds on my system.

Using Selected Images

When you find the image you want, you open it by clicking the appropriate button. The image is then viewed in a window which can be zoomed to the size of the Macintosh window. Scroll bars are

provided for precise (but slow) positioning of the image. No grabber hand is available for faster positioning.

The image may be viewed, cut or copied, printed, or saved in a different format. The cutting and copying is done by choosing **Select All** from the **Edit Menu** or by using the selection rectangle. The printing involves use of standard **Page Setup** and **Print** dialog boxes. You may also view or change the key words associated with the displayed image or may view the list of all key words in the catalog.

Converting Image Formats

The most useful and straightforward feature of Curator is the ability to convert images from one format to another. This is accomplished merely by checking the desired new format, then using the SAVE AS... command. The supported conversions are indicated by dimming the box for those that are not supported (e.g., MacPaint to Mac EPS.)

The Curator Manual

The manual says, "You may not need to read this manual at all!" That may be true, but only if you are the one that wrote Curator. This program is about as intuitive as driving a car in Boston. After a lot of driving and many wrong turns, maybe you will get there. The manual usually (but not always) does an adequate job of explaining the product. The section of **Browsing** is confusing. It seems to contain an irrelevant illustration (i.e., it shows the wrong window) and leaves out a critical step.

The Bottom Line

Solution International, Inc., had good intentions in developing Curator. This program addresses a problem which needs to be solved (i.e., the multiplicity of graphics formats.) The implementation, however, is in my opinion, not up to this company's usual standards. The interface of the art management part of the program needs to be simplified. Also, the manual needs to be clarified, especially as it describes the **Browsing** function.

The most important function of Curator, however, is image conversion and this function is simple and powerful. It does what it claims it will do.

If you need the ability to convert graphics from one format to another, buy Curator. If all you need is an art management program, wait until it is refined.

EOF

FOR THOSE WHO NEED TO KNOW 68 MICRO JOURNAL™

Pascal

A Tutorial

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This month we are going to look at variable allocation, but before we do, let's take a look at "structured" constants. Structured constants refer to constant arrays and/or records. These tend to straddle the line between constants and variables in all three languages. Pascal and Modula-2K both have true structured constants. C uses an equally useful mechanism of initialized variables.

In Pascal, a structured constant in defined in one of three ways :

```
const
  abc = local <type> <constant-list> ;
  abc = entry <type> <constant-list> ;
  abc = external <type> ;
```

A local structured constant is only known within the module or procedure where it is declared. An entry structured constant in known within the module where it is defined, and any other modules that declare the same name as external. Entry/external structured constants cannot be defined within a procedure. The constant list consists of the constant value in the case of simple types, such as integer, boolean, string, or constants within in parenthesis in the case of arrays or records (one level of parenthesis for each level of record or array). For instance, all of these are valid :

```
const
  abc = local integer 0 ;
  def = entry string 'this is a test' ;
```

```
  ghi = local array [1..5] of integer
(10,100,1000,50,70) ;
  jkl = entry record
        a : integer ;
        b : string ;
        c : array [1..3,1..4] of integer
      end
(25,'test',((5,10,15,20),(10,20,30,40),
(20,40,60,80))) ;
```

In Modula-2K the distinction if the constant is local, entry, or external is determined by whether or not the constant is mentioned in the definition module (indicating export (entry)), or whether it is in an import list (indicating import (external)). If it is in neither of these, then it is a local structured constant. As in the Pascal, if it is external (imported), then the constant-list is not included. Also, the constant list is not specified when defining an exported constant in the definition module. So for Modula-2K :

DEFINITION MODULE

```
                                    {local}
const abc = <type>                  {entry}
                                    {external}
```

IMPLEMENTATION MODULE

```
const abc = <type> <list>    {local}
const abc = <type> <list>    {global}
```

from <module> import abc   {external}

In Pascal and Modula-2K these constants can be indexed, and fields accessed just like a structured variable.

In C a totally different concept is used, that of the initialized variable. Whether or not the variable is local to the module follows the same rules as normal variables. The method of definition is almost the same as Pascal and Modula-2K :

```
int i = 0 ;
int x[4] = {10, 20, 30, 40} ;
```

The implementation of initialized variables varies depending on the operating system environment, but in general the data is copied from the memory module to the data area before execution of the C program, resulting in two copies of the data in memory when the program starts.

In the Pascal and Modula-2K implementation, the data stays in the memory module, since it is not to be modified. Wanting to have an actual initialized variable (versus a set of constants, such as a look-up table) is rare, but when required can be done with just a simple assignment, resulting in essentially the same situation as in C's initialized variables.

Now onto real variables. The default allocation method in Pascal is on the stack. This applies to any variables declared globally, or local to a procedure. This is the only allocation method in standard pascal. In a modular environment it is vital that the global stack variables be declared identically in all modules, this is assured by declaring them in an include file. Modula-2K and C do not have a global stack, but have a local stack for local procedure variables.

"Register" and "FPR" variables are also allowed. Register variables are stored in the CPU registers (D4-D7) and up to 4 are allowed per procedure. Likewise, FPR variables are stored in the 68881 registers (FP2-FP7) and up to 6 are allowed per procedure. Both of these types are available in Modula-2K as well. C allows register variables, how many are allowed per procedure, and whether or not they

are stored in the CPU or FPU registers would naturally depend on the specific compiler used.

PASCAL	MODULA-2K

```
i : integer register    i [register] : integer
```

```
                      C

              register int i;
```

The "Varib" type is unique to Pascal, since it's default type is stack storage, and the other two languages default to the equivalent of "Varib". "Varib" simply means that the storage location of the variable (relative to a base register) is determined by the linker, not the compiler. In Pascal "Varib" may store up to 32K bytes of data. In Modula-2K and at least one C compiler, up to 64K bytes of data is allowed. In Pascal, Modula-2K and at least one C compiler, another section called "Remote" is allowed, which is like "Varib", but uses 32 bit address and allows up to 2 giga bytes of storage.

There are also variants of these type depending on whether you want the variable to be known outside of this module. If the variable is declared inside a procedure, then there are no possible options, it is not known outside of that procedure (normal scope rules).

In Pascal if there is no special modifier after the "Varib" or "Remote" then it is not known outside of the module or procedure where it is declared. To make the variable known outside of this module, it is followed by the word "Entry", and to access this variable in another module, by the word "External".

```
i : integer varib {local}
i : integer varib entry {global}
i : integer varib external {global, defined
                                  elsewhere}
i : integer remote entry {remote global}
```

In Modula-2K the distinction is defined by the appearance in the definition module or an import list :

DEFINITION MODULE

```
                              {local}
i : integer                   {global}
                              {global, defined
                               elsewhere}
i [remote] : integer          {global remote}
```

IMPLEMENTATION MODULE

```
i : integer                   {local}
i : integer                   {global}
from <module> import i        {global, defined
                               elsewhere}
i [remote] : integer          {global remote}
```

In C a variable declared outside of a procedure without any modification of the declaration type is essentially a globally known variable (the same as varib entry in Pascal), and can be referenced in other modules by using the keyword "extern" (the same as varib external in Pascal). To make a variable that is global to a module, but not known outside the module, the keyword "static" is used.

```
static int i {local}
int i {global}
extern int i {global, defined elsewhere}
remote int i {global remote}
```

One very useful mode found in Pascal and Modula-2K is the absolute addressing mode. This allows a variable to be located at a specific location in the memory map, which is mostly used for I/O ports. As examples :

PASCAL

```
i : integer at $ff104 ; {I/O port}
a : array [1..24,1..80] of char at $e0000 ;
                        {memory mapped screen}
```

MODULA-2K

```
i ($ff104) : integer ; {I/O port}
a ($e0000) : array [1..24,1..80] of char ;
                      {memory mapped screen}
```

In C, this facility is not available, and pointers are often used to access absolute addressing by assigning the address to the pointer variable. This method unfortunately generates more code, but there is a way of accessing absolute addresses using the same code as Pascal and Modula-2K by using a "caste" :

```
#define i * (integer *) 0xff104
```

Unfortunately this does not seem to work for arrays or records, and so you are back to using pointers or defining each individual field of a record at a specific address.

Next time we will look at operators.

EOF.

FOR THOSE WHO NEED TO KNOW   **68 MICRO JOURNAL**™

A condensed outline of utilites that many have asked for. They are now available from

South East Media (see catalog.)

GCS FILE TRANSFER UTILITIES

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GCS FILE TRANSFER UTILITIES are a group of programs to transfer files on MS-DOS (PC) and FLEX format floppy disks to and from OS-9.

MS-DOS text files can be transferred, but not binary files. Remember that Intel and Motorola operation codes are incompatible. Both text and binary FLEX files can be transferred. Binary FLEX files are not directly executable by OS-9.

The use and syntax of these utilities is standard OS-9. Usually, the default values for number of sides, sectors, etc., will suffice. A number of options are available so that other standard MS-DOS and FLEX formats, as well as some other formats, can be read and written.

NOTES - General:

Options can be displayed by entering: <prognam> -?.

Optional <devnam> for the device where the PC or FLEX disk is located must begin with a '/', (for example: /D1).

In the program descriptions which follow, the default values are assumed.

Certain features of OS-9 and MS-DOS file handling such as complete hierarchical handling, have not been included in these programs.

First level MS-DOS sub-directories can be read by PCDIR, PCREAD, PCDELETE and PCRENAME.

NOTES - 40 and 80 Track Drives:

MS-DOS files are normally written on disks formatted for 40 tracks on a 40 track drive. This is generally true for FLEX format disks. Assuming that your drives are properly aligned, you will be able to read MS-DOS and FLEX disks on either 40 or 80 track drives. The defaults are set for 80 track drives: use the -4 option for a 40 track drive. (For EXPLORE, the sense is reversed - the default is 40 track and the option is -8). Writing files to MS-DOS or FLEX disks should be done using a 40 track drive, if the disks were originally formatted on 40 track drives. If an eighty track drive is used to write to disks formatted on a 40 track drive, it is unlikely that any files so written will be read reliably on a PC or FLEX system. See the GMX Micro-20 Manual Addenda for OS-9 (page 3-2).

PCDIR

This program displays the directory of a PC format disk in /D0.

The program will read the root directory or a directory in the root directory. The path for a directory in the root directory is: <dir>.

Usage

Function: display pc format disk file directory

Syntax: pcdir [<opts>] [<devnam>] [<opts>]

Options: -e extended directory listing -f display disk format -d single density -s single sided -4 40 track drive -c=<num> sectors per cluster -n=<num> sectors per track -t=<num> tracks per surface

-? usage

Default - options: 80 track drive, double density, double sided, 2 sectors per cluster, 9 sectors per track, 40 track disk, 512 bytes per sector

- devnam: /d0

PCDUMP

This program displays individual sectors of a PC format disk in /D0.

The user is prompted for track, sector and side.

After a sector has been displayed, the user is prompted for:

N - next sector P - previous sector R - repeat current sector S - start over Q - quit

Output is to <stdout>. Redirection can be used for output to a file or printer.

Usage

Function: display pc format disk files by specified track, sector and side

Syntax: pcdump [<opts>] [<devnam>] [<opts>]

Options: -d single density -s single sided -4 40 track drive -c=<num> sectors per cluster -n=<num> sectors per track -t=<num> tracks per surface

-? usage

Default - options: 80 track drive, double density, double sided, 2 sectors per cluster, 9 sectors per track, 40 track disk, 512 bytes per sector

- devnam: /d0

PCREAD

This program reads a file contained in a PC format disk in /D0.

The file must either be in the root directory or a first level sub-directory.

The path for a file in the first level sub-directory is:

<dir><file>.

Line feeds are stripped from output. Use the <-l> option to retain the line feeds.

Output is to <stdout>. Redirection can be used for output to a file or printer.

Usage

Function: read pc format disk files

Syntax: pcread [<opts>] [[devnam] <path> [dev-nam] [<opts>]

Options: -l pass line feeds -d single density -s single sided -4 40 track drive -c=<num> sectors per cluster -n=<num> sectors per track -t=<num> tracks per surface

-? usage

Default - options: 80 track drive, double density, double sided, 2 sectors per cluster, 9 sectors per track, 40 track disk, 512 bytes per sector, strip line feeds

- devnam: /d0

PCWRITE

This program writes a OS-9 file to a file on a PC format disk in /D0.

The file may be written only into the root directory.

If no destination filename (dstfile) is specified in the command line, the OS-9 file (srcfile) is used as the destination filename).

If the OS-9 filename is longer than 8 characters, the destination filename written into the PC disk will be truncated. If the desti- nation filename is included in the command line, a more suitable filename and extension can be used (8 characters for filename - 3 characters for extension).

NOTE: Please observe the caution on writing to 40 and 80 track drives. For reliable results, 40 track disks should only be written with 40 track drives. Usage

Function: write pc format disk files

Syntax: pcwrite [<opts>] [<devnam>] <srcfile> [<dstfile>] [<devnam>] [<opts>]

Options: -l don't add line feeds -d single density -
s single sided -4 40 track drive -c=<num> sectors per
cluster -n=<num> sectors per track -t=<num> tracks
per surface

-? usage

Default - options: 80 track drive, double density,
double sided, 2 sectors per cluster, 9 sectors per track,
40 track disk, 512 bytes per sector, add line feeds

- devnam: /d0

PCDELETE

This program deletes a file on pc format disk.
Please be careful with this program! Once a file is
deleted, it cannot be restored because the cluster chain
for the file is cleared from the file allocation tables.
This is done so that the space on the disk can be reused
to write other files.

The prompt "pcdelete: Deleting: <filename> - Are
you sure (Y/N) ?" provides an opportunity to abort the
delete.

The path for a file in a first level sub-directory is:

<dir><file>.

Usage

Function: delete pc format disk file

Syntax: pcdelete [<opts>] [devnam] <filename>
[devnam] [<opts>]

Options:

-d single density -s single sided -4 40 track drive
-c=<num> sectors per cluster -n=<num> sectors per
track -t=<num> tracks per surface

-? usage

Default - options: 80 track drive, double density,
double sided, 2 sectors per cluster, 9 sectors per track,
40 track disk, 512 bytes per sector

- devnam: /d0

PCRENAME

This program renames a file contained on a pc
format disk.

The path for a file in a first level sub-directory is:

<dir><file>

Usage

Function: rename a pc format disk file

Syntax: pcrename [<opts>] [<devnam>] <old file
name> <new file name>

Options:

-d single density -s single sided -4 40 track drive
-c=<num> sectors per cluster -n=<num> sectors per
track -t=<num> tracks per surface

-? usage

Default - options: 80 track drive, double density,
double sided, 2 sectors per cluster, 9 sectors per track,
40 track disk, 512 bytes per sector

- devnam: /d0

PCFORMAT

This program formats a floppy disk with pc format.
While a disk can be formatted on either a 40 or 80
track drive, the disk is actually formatted as 40 track. If
a disk is formatted on a 80 track drive and is always
written on a 80 track drive, the disk can be read on
either a 40 or 80 track drive.

For reliable results, format and write on 40 track
drives, since pc systems are normally 40 track.

It is not normally necessary to verify formatting
given the quality of most current day floppy disks. The
-q option verifies the first sector on each track. The -v
option verifies all sectors on all tracks. Verifying does
slow down the format process.

Usage

Function: format pc disk

Syntax: pcformat [<opts>] [<devnam>] [<opts>]

Options:

-l prompts for volume label -q quick verify -s single sided -v verify -4 40 track drive

-? usage

Defaults - options: 80 track drive, double density, double sided, 40 track disk, 9 sectors per track, 2 sectors per cluster, 512 bytes per sector, no volume label, no verify

devnam: /d0

FLEXDIR

This program displays the directory of a FLEX format disk in /D0.

Usage

Function: display FLEX format disk file directory

Syntax: flexdir [<opts>] [<devnam>] [<opts>]

Options: -e extended directory listing -s double sided -n=<num> sectors per track

-? usage

Default - options: single density, single sided, 0/ 80 track disk, 10 sectors per track, 0/80 track drive - devnam: /d0

FLEXDUMP

This program displays individual sectors of a FLEX format disk in /D0.

The user is prompted for track, sector and side.

After a sector has been displayed, the user is prompted for:

N - next sector P - previous sector S - start over Q - quit

Output is to <stdout>. Redirection can be used for output to a file or printer.

Usage

Function: display specified FLEX format disk file sector

Syntax: flexdump [<opts>] [<devnam>] <file> [<devnam>] [<opts>]

Options: -d double density -s double sided -4 40 track drive -n=<num> sectors per track -t=<num> tracks per surface

-? usage

Default - options: 80 track drive, single density, single sided, 40 track disk, 10 sectors per track,

- devnam: /d0

FLEXREAD

This program reads a file contained in a FLEX format disk in /D0.

Use the -b option to transfer binary files directly to a OS-9 file.

If formatted binary output is desired for screen or printer output; also use the -f option.

Usage

Function: read FLEX format disk files

Syntax: flexread [<opts>][<devnam>] <file> [<devnam>][<opts>]

Options: -b binary file -f format binary output -d double density -s double sided -4 40 track drive - n=<num> sectors per track -t=<num> tracks per surface -? usage

Default - options: 80 track drive, single density, single sided, 40 track disk, 10 sectors per track, text file

- devnam: /d0

FLEXWRITE

This program writes a OS-9 file to a file on a FLEX format disk in /D0.

If no destination filename (dstfile) is specified in the command line, the OS-9 file (srcfile) is used as the destination filename.

If the OS-9 filename is longer than 8 characters, the filename (dstfile) written into the FLEX disk will

be truncated. If the desti- nation filename is included in the command line, a more suitable filename and extension can be used (8 characters for file - 3 characters for extension).

Embedded tab characters (09) in the source file file are replaced by a tab character followed by a tab count character. The default tab count character default if 4. This can be changed use of -e option.

Strings of space characters ($20) in the source file are replaced by a tab character ($09) followed by a space count character. This creates normal FLEX space compression.

NOTE: Please observe the caution on writing to 40 and 80 track drives. For reliable results, 40 track disks should only be written with 40 track drives.

Usage

Function: write FLEX format disk files

Syntax: flexwrite |<opts>| |<devnam>| <srcfile> |<dstfile>| |<devnam>| |<opts>|

Options: -b binary file -e=<num> spaces per tab -d double density -s double sided -4 40 track drive -n=<num> sectors per track -? usage

Default - options: text file, 4 spaces per tab, 80 track drive, single density, single sided, 40 track disk, 10 sectors per track, 256 bytes per sector

 - devnam: /d0

EXPLORE

This program can be used to display files on a variety of disk file formats.

It should work on disks which have been formatted using one of the IBM formatting algorithms with a controller using a WD1770, WD1772, WD1773 or WD179x chip.

Usage

Function: display disk files by specified track, sector and side

Syntax: explore [<opts>] [<devnam>] [<opts>]

Options: -g GMX Micro-20 configuration -b 512 bytes per sector -d double density - 16 sectors per track -s double sided -q 80 track density disk -8 80 track drive -n=<num> sectors per track -t track 0 always single density -m strip most significant bit of data byte -? usage

Default - options: 40 track drive, single density, single sided, 10 sectors per track, all tracks same density, 40 track disk, 256 bytes per sector, pass most significent bit of data byte

 - devnam: /d0

INSTALLATION and TUTORIAL

This is a simple tutorial for the GCS File Transfer Utilities package.

First copy the utilities into your CMDS directory using the "load_util " procedure file.

Enter the following command line:

 pcformat -? This will display the syntax and options for this command. (-? has the same function for other commands.)

Insert an unformatted disk into D0 device - it is assumed that this is an 80 track drive. If it is a 40 track drive, be sure to use the -4 option.

If you want to use another device; use /D1, for example, on the command line.

Enter the following command line:

pcformat

The program will display the device on which the disk to be formatted is mounted (/D0 for default). You will be asked if you wish to continue, enter: y . Formatting will proceed; at the end of a successful format, a summary of good and bad sectors will appear. If formatting should abort, an explanatory message will appear.

Use the CHD command to set up the OS-9 source/ destination directory.

Now try writing a file to the MS-DOS disk.

Enter the following command line:

pcwrite file1

(file1 = any text file in the current data directory)

Note that the file name may be truncated if it is longer than 8 characters since that is the maximum length for a MS-DOS file name.

Check that this has taken place by entering the following command line:

pcdir 'file1' will be the only file shown in the directory.

Read this file disk by entering the following command line:

pcread file1

This will display 'file1' on your terminal. Note that extraneous line feeds have been stripped from the display. If you want to keep line feeds, use the -l option. On the terminal, this option should result in double spacing.

Enter the following command line:

pcread file1 >-another_name This command line will write 'file1' into the current data directory. Don't forget the '-' after the '>' and use a different file name.

Check that this has taken place by entering the following command line:

list another_name

Usually the default values will work for most MS-DOS disks, since this is the most common current format. The 'pcdir' command can read several standard MS-DOS formats. The -f option for this command will indicate which format is used so that you can use appropriate options with the other commands if necessary.

Enter the following command line:

pcdir -? This will display the syntax and options for this command.

Enter the following command line:

pcdir

This will display the files in the directory of the MS-DOS disk - just showing the names of the files.

Enter the following command line:

pcdir -e This will display an extended listing of the directory.

Enter the following command line:

pcdir -f This will display the format parameters of the MS-DOS disk if it is one of the standard formats.

Using the preceeding command lines, you will receive an idea of the use of this command. Options can be combined as needed. They can go before or after filenames. Likewise device name.

As you can see these are a very complete set of utilities. While there have
been other utilities available to our readers. This seems to be the most complete set we have seen so far.

Available for 68XXX OS-9 and CoCo OS-9 from S.E. MEDIA.

Price: $99.00

EOF

Bit-Bucket

By: All of us

"Contribute Nothing · Expect Nothing", DMW '86

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FAX 515-224-1352

FOR MORE INFORMATION CONTACT:
Mr. Andy Ball
Vice President, Marketing
Microware Systems Corporation
1900 NW 114th Street
Des Moines, Iowa 50322
515-224-1929

MICROWARE ANNOUNCES THE RELEASE OF DEVELOPMENT PAKS FOR THE MOTOROLA VME 133A, 134 & 135

DES MOINES, Iowa. — Microware Systems Corporation announces the release of three new OS-9 Operating System Development Paks for the Motorola VME 133A, 134 and 135 MC68020 monoboard microcomputers. OS-9 Development Paks provide a complete C Language development environment including a Kernighan & Ritchie compatible C compiler, assembler, linker, debugger, screen editor and over 50 other popular utility commands. Microware's OS-9 real-time operating system is used by major manufacturers world wide in a variety of applications ranging from embedded process control systems, multi-user development systems and consumer electronics.

These OS-9 Development Paks have been customized to support a variety of controllers for floppy and hard disks, magnetic tape, and serial and parallel I/O. Controllers supported by Microware's Development Paks include the MVME 319, 320, 335, 350 and 050.

The Microware MVME 133A Development Pak includes complete device driver or trap handler support for the on-board MC68881 FPU, RS-232 serial ports and 20MHz clock. These features make this microcomputer a popular choice for time-critical or numerically intensive industrial applications.

The MVME 134 monoboard microcomputer includes a MC68851 Paged Memory Management Unit (PMMU). OS-9 memory protection support for this device and 4 Mbytes on-board RAM make the MVME 134 Development Pak an ideal engine for large multi-user development systems.

A significant feature of the MVME 135 is a VSB private memory bus. The VSB bus optimizes VME bandwidth utilization by minimizing CPU accesses to the VME-bus. OS-9 support for this private memory bus makes the MVME 135 especially useful for multi-processor systems or systems incorporating high-speed DMA controllers.

The three OS-9 Development Paks are immediately available from Microware and authorized Microware distributors. Please contact Microware for pricing information.

The OS-9 Operating system is a real-time, multi-user and multi-tasking system for computers based on the Motorola family of 68xxx processors. OS-9 is compact, ROMable and provides a UNIX-style environment for application software. Since its introduction in 1983, OS-9/68000 has been licensed to over 350 original equipment manufacturers (OEMs) world-wide for use in a variety of industrial, scientific and consumer products.

Computer Systems Consultants, Inc.
1454 Latta Lane
Conyers, GA 30207

Don Williams, Editor
68 Micro Journal
5900 Cassandra Smith
Hixson, TN 37343

Dear Don:

As Leo Taylor mentioned in his letter to the editor in 68 Micro, the current version of Super Sleuth for the 68010 was designed to disassemble 68000, 68008, and 68010 machine code, which it does. This current version has no current known bugs. I have used it for some substantial disassemblies, and customers have reported successfully using it for disassembling several large programs, including monitors and operating systems.

The OS9/68000 version of Super Sleuth does not currently process the OS9/68000 header and system calls; however, it will process the remainder of the machine code properly. The table of system calls in Super Sleuth is not currently used, but is for use in the future version. Both the format of the OS9/68000 file header and the exact format of the OS9/68000 system calls are badly-documented, making this task much more difficult than it was under OS9/6809.

This system-specific information will be processed in a future version of Super Sleuth which will also process the UNIX V header and system calls.

Incidentally, Leo Taylor's description of the UNIX file system is not accurate. It is substantially more robust then his description would indicate. There are also two major current versions of the file system, as originated by Berkeley and as originated by AT&T. Those who are interested may read one of the many books describing the internals of the UNIX operating system.

Thank You,

E. M. (Bud) Pass

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Dear Don,

How time flies! Hadn't realised it's been so long since I began writing about logic functions in BASIC, as I've been so busy with my 68000 version of RBASIC, plus a design project which suddenly cropped up. So let's see if I can pick up the thread of where I left off.

I was emphasising just how careful we have to be when modifying logic expressions which include ELSE as part of their evaluation. Now for a much different example, which begins by looking like a mere extension of my earlier discussion ... but just wait till the punch-line! Let's look at the following :

```
100   IF X% < 3 GOTO 120
110   IF Y% = 9 GOTO 130
120   IF X% < 2 THEN Z% = 3
130   rest of program
```

Note that we can arrive at Line-120 by two different routes from Line-100 (i) directly if (X%<3) or (ii) failing that, and falling through to Line-110, if (Y% <> 9). So let's try combining Lines 100 and 110 into

```
100   IF X% > 2 AND Y% = 9 GOTO 130
120   IF X% < 2 THEN Z% = 3
130   rest of program
```

This is permissible because the program will arrive at Line-110 in the original program only if (X% > 2), and then only if (Y% = 9) will it bypass Line-120 and GOTO 130. We haven't changed the intent of the original at all, as the program will still arrive at Line-120 if (X% < 3) or if (Y% <> 9). Aha! So how about now combining Lines 100 and 120 so

```
100  IF (X% < 3 OR Y% <> 9) AND X% < 2 THEN Z% = 3
130  rest of program
```

I hope you'll remember from my 'Logically Speaking' series how to negate the two conditionals of the old Line-100 to produce the OR-expression inside the parens in our new Line-100. If not, I'll summarise it as "The opposite of (X% > 2) is (X% < 3); the opposite of AND is OR; and the opposite of (Y% = 9) is (Y% <> 9). So our old program will reach Line-120 under this combination, ie, (X% < 3 OR Y% <> 9), and then if (X% < 2) Z% will get set to 3. Don't be tempted to omit the parens around the pair of OR-coupled terms, otherwise BASIC will read the line as

```
100  IF X% < 3 OR (Y% <> 9 AND X% < 2) THEN Z% = 3.
```

Remember that AND equates to multiplication, and OR to addition, so the two terms coupled by AND will have a higher mathematical priority than OR, and BASIC will see a set of implied parens around the AND-coupled terms.

Normally this would be about as far as we could go with reducing the complexity of the original program, but now we have a whole army of Boolean Algebra experts out there as a result of my 'Logically Speaking' tutorial series, and I doubt that they'll be content to let things rest there. As a first step, they'd expand the logic terms into

```
(X% < 2).(X% < 3) + (X% < 2).(Y% <> 9)
```

and then they'd observe that the first product-term can be reduced to a simple (X% < 2). After all, if something is less than 3 and at the same time is less than 2, the condition can obviously be met by the simple condition less than 2. So now we have

```
(X% < 2) + (X% < 2).(Y% <> 9)
```

And, as a final step, we'll apply Rule 5 of the Laws of Boolean Algebra (see page 23 of the September 1987 issue of 68%J), which states that 'where a term in its entirety forms part of a larger term, the larger term can disappear.' This, of course - where (X% < 2) corresponds to 'a' and (Y% <> 9) to 'b' - reduces us to a mere

```
(X% < 2)
```

so we can now simplify our program to

```
100  IF X% < 2 THEN Z% = 3
130  rest of program
```

All of which, though it's not intuitively obvious in our original program, tells us that the original Lines 100 and 110 are redundant, their only function being to act as 'program-complicators' and 'time-wasters'.

I hope this little example will serve to encourage others to apply Boolean Algebra to their BASIC programs, as well as to circuit design!!

And now, before signing off, I have to confess to having fallen into a trap of my own making. I should have known better, but a reader, Dexter S. French, pointed out that in my number-base conversion program (in XBASIC XPLANATIONS) I said that the following

```
40  K%=ASC(MID$(X$,I%,1))-48: IF K%>9 THEN K%=K%-7
```

could be shortened, by using logic expressions, into

```
40  K%=ASC(MID$(X$,I%,1))-48 +7*(K%>9)
```

which, I'm afraid, just ain't true! This comes of not checking out one's own program - - no matter how straightforward it looks. The first Line-40 is OK, because by the time the conditional (K%>9) is examined, it has already been calculated in the first part of the line.

Unlike pocket-calculators, however, BASIC doesn't update the value of a variable until it's completed its whole chain of calculations in a set of Math-registers, so the second Line-40, in the latter portion, is basing its logic decision on a previous value of K%, and not on its value at the point where it's just subtracted 48. The math operation is incomplete at this point, as there's still something to be added to the continuing evaluation of K%, so K% is not yet due to be updated.

Although I'm thoroughly aware of the fact that one shouldn't include logic-functions where the value of the variable has not yet been determined, I slipped up, and this one sort of got by me. Sorry about that! Please amend your copies of my 'XBASIC XPLANATIONS' accordingly.

I'll be in touch.

Sincerely,

Don Williams,
68 Micro Journal,
5900 Cassandra Smith Road,
Hixson, TN 37343

R. Jones
President

MOTOROLA ANNOUNCES 32-BIT IEEE 754-1985 COMPLIANT FLOATING POINT DIGITAL SIGNAL PROCESSORS

While fully supporting the IEEE 754-1985 binary floating point standard, the DSP96001 general purpose digital signal processor has a peak performance of 40 million floating point operations per second (MFLOPS) using a 26.7 MHz oscillator

Motorola Inc., Austin, Texas, March 30, 1988.......Motorola's Digital Signal Processor Operation, based in Austin, Texas, producer of the HYPERformance (TM) DSP56000 family of 24-bit digital signal processors, is announcing plans to launch an extension of its existing family. The DSP96001 will be the first available offering in the new DSP96000 Family equipped with floating point.

The DSP96001 is a 32-bit x 32-bit floating point digital signal processor that has 96-bit accumulators. The DSP96001 is in total conformance with the IEEE 754-1985 binary floating point standard. Software for the DSP96001 is both upward and downward compatible with the software for the DSP56000 family of fixed point digital signal processors. There are 512 words of on-chip full speed program RAM (PRAM), two 512-word data RAMS, two pre-programmed data ROMs, special on-chip bootstrap hardware for efficient program loading into the PRAM, an on-chip debug circuit to allow access to internal resources in support of the On-Chip Emulation, OnCE(TM), and two full DMA channels.

According to Bryant Wilder, Motorola's Manager for the digital signal processor operation, "the in-circuit emulation characteristics of board level products has been carried one step further. Motorola has brought that degree of debugging capability on-chip with the OnCE (TM) feature of the DSP96000 Family."

There are three 32-bit execution units operating in parallel within the CPU. These are the Data ALU, the Address Generation Unit, and the Program Controller. Having an MCU-like architecture with on-chip peripherals, program and data memory, and expansion ports, the DSP96001 offers substantial flexibility to the system designer for a variety of digital applications. Its programming model and instruction set are similar to that of a microprocessor, making the generation of efficient, compact code, straightforward for the programmer.

Operating at a peak performance rate of 40 MFLOPS, the DSP96001 is an ideal candidate to support a wide variety of real time DSP applications that demand the accuracy of an IEEE conformant floating point unit. Some of the more familiar applications would include high-speed controllers, digital audio systems, numeric processing, image and speech processing, spectral analysis, instrumentation, medical environments, and navigational systems.

Features that make the throughput possible for these application areas include the following:

Speed -- In less than 2 milliseconds, a 1024 point complex floating point Fast Fourier Transform (FFT) can be executed by the DSP96001 operating at a 13.33 million instructions per second (MIPS) rate.

Precision -- In single precision (SP), there is full IEEE 754-1985 floating point conformance. Full 32-bit buses and ALU support the SP operations by performing all floating point operations in single extended precision (SEP) arithmetic. This allows the SP error bounds to be met after only one of the four IEEE rounding operations have been executed - round to +infinity, - infinity, zero, and even. The general register file consists of ten 96-bit registers. IEEE double extended precision operations performed in software are supported. Enhancing precision one step further, the data ALU also supports 32-bit integer arithmetic including 32 x 32 integer multiplication with a full 64-bit product.

Parallelism -- The Data ALU, Address Generation Unit, and Program controller operate in parallel within the CPU so that an instruction prefetch, up to three floating point operations, two data moves, and two address pointer updates using one of three types of arithmetic (linear, modulo, or reverse carry) can be executed in a single instruction cycle. The 40 MFLOP peak performance is made possible as a result of this parallelism. Two on-chip DMA controllers operate unobtrusively in parallel with the CPU to assure there is an adequate supply of data available for processing.

Integration -- The DSP96001 is highly integrated. In addition to the three independent execution units and the two DMA controllers, there are six on-chip memories, three on-chip MCU style peripherals (serial communications interface, synchronous serial interface, and a 32-bit host interface), a clock generator and eight 32-bit wide buses (three address and five data). This makes the overall system high performance, low power, and compact -- an excellent cost performance value.

DSP96001 Compatibility -- The DSP96001 has the same basic architecture as the DSP56000 Family. Its instruction set is a superset of the DSP56000 Family instruction set. This compatibility means the hardware and software development tools are nearly identical for all Motorola DSPs, and therefore, investment in existing software is preserved. Applications can be easily developed on one processor and migrated across the Motorola DSP product line to meet various cost and performance objectives.

ATARI ST1040

Geza Holzhaker,
Apdo. Correos # 393,
Merida 5101, VENEZUELA.

Dear Mr. Williams:

I've just read the "ATARI call" in the Jan.88 issue.
I'd like very much if the 68MJ would include a regular section for the Atari computer.
Until today I've bought two ST1040's (Monochr. res.1 640x400). When the ST1040 came out, it had the worst Basic I've seen in my life. But now there is available GFA_Basic, it's structured and very fast, it has graphic capabilities and easy access to 68000 calls. There is also a useful assembler/disassembler (Assempro) to help writing 68000 codes, PC Board design, Publishing Partner and more.
I think the ST1040 is not expensive taking into account the graphic capabilities that other machines are lacking. The keyboard is not comfortable, but with all the software appearing the ST1040 is definitively not a toy.
I wish I could buy OS9 for the Atari some day, I am sure the Basic09 and "C" from Microware are better than others, but I don't know if they support graphics. I hope SK*DOS would be released too.

I've already installed a 5" (80 tracks) drive as drive "B" instead of another 3" drive. It works perfectly, even without slowing down the stepping rate of the drive (Teac-55F1. I feel 5" disks more reliable than 3" disks. I installed the 5" drive without having the schematics of the ST1040. I'd appreciate very (but very) much if somebody could send me a copy (and related info) of it to me. It is impossible to get them here in Venezuela.

At the present time we are developing Intelligent Electrocardiographic and Electroencephalographic systems with the ST1040.

I hope the "ATARI" section will be a regular issue in the 68MJ.

Sincerely yours,

GMX MICRO-20 and TWINGLE-20 PRICE LIST

All versions include 1 SAB Board

	MICRO-20 with 1MB RAM	MICRO-20 with 2MB RAM	TWINGLE-20 with 4MB RAM
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OPTIONAL PARTS AND ACCESSORIES

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MOTOROLA 68020 USERS MANUAL	$ 18.00
MOTOROLA 68030 USERS MANUAL	$ 18.00
MOTOROLA 68881 USERS MANUAL	$ 18.00

SBC ACCESSORY PACKAGE (M20-AP) $1399.00

The package includes a PC-style cabinet with a custom backpanel, a 25 Megabyte (unformatted) hard disk and controller, a floppy disk drive, a 150 watt power supply, cooling fan, panel mounted reset and abort switches, and all necessary internal cabling. (For use with SAB—9D serial connectors only.)

2nd 5"80 FLOPPY & CABLES FOR M20-AP, AOD	$ 250.00
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TO SUBSTITUTE 50MB HD FOR 25MB HD, AOD	$ 290.00
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I/O EXPANSION BOARDS

16 PORT SERIAL BOARD ONLY (SBC-16S) ....$ 335.00

The SBC-16S extends the I/O capabilities of the GMX Micro-20 68020 Single-board Computer by adding sixteen asynchronous serial I/O ports. By using two SBC-16S boards, a total of thirty-six serial ports are possible.

RS232 ADAPTER (SAB-25, SAB-9D or SAB-8M) $165.00

The board provides level-shifting between TTL level and standard RS-232 signal levels for up to 4 serial I/O ports.

60 LINE PARALLEL I/O BOARD (SBC-60P) $398.00

The GMX SBC-60P uses three 68230 Parallel Interface/Timers (PI/Ts) to provide up to forty-eight parallel I/O lines. The I/O lines are buffered in six groups of eight lines each, with separate buffer direction control for each group. Buffer direction can be fixed by hardware jumpers, or can be software programmable for bidirectional applications.

PROTOTYPING BOARD (SBC-WW) $75.00

The SBC-WW provides a means of developing and testing custom I/O interface designs for the GMX Micro-20 68020 Single-board Computer. The board provides areas for both DIP (Dual Inline Package) and PGA (Pin Grid Array) devices, and a pre-wired memory area for up to 512K bytes of dynamic RAM.

I/O BUS ADAPTER (SBC-BA) $195.00

The SBC-BA provides an interface between the GMX Micro-20 68020 Single-board Computer and the Motorola Input/Output Channel (I/O bus). With the I/O bus, up to sixteen off-the-shelf or custom peripheral devices (I/O modules) can be connected to the GMX Micro-20.

ARCNET LAN board w/o Software (SBC-AN) $475.00

The SBC-AN provides an interface between the GMX Micro-20 68020 Single-board Computer and the ARCNET modified token-passing Local Area Network (LAN) originally developed by Datapoint Corp. The ARCNET is a baseband network with a data transmission rate of 2.5 Megabits/second. The standard transmission media is a single 93 ohm RG-62/U coaxial cable. Fiber optic versions are available as an option.

OS9 LAN Software Drivers for SBC-AN 120.00

GMX MICRO-20 SOFTWARE

020 BUG UPDATE — PROMS & MANUAL $150.00

THESE 68020 OPERATING SYSTEMS ARE PRICED WHEN PURCHASED WITH THE MICRO-20. PLEASE ADD $150.00 IF PURCHASED LATER FOR THE UPDATED PROMS AND MANUALS. ALL SHIPPED STANDARD ON 5¼" DISKS 3½" OPTIONAL IF SPECIFIED.

OS9/6 020 PROFESSIONAL PAK $850.00
Includes O.S., "C", uMACS EDITOR, ASSEMBLER, DEBUGGER, development utilities, 68881 support.

OS9/6 020 PERSONAL PAK ....$ 400.00
Personal OS-9 systems require a GMX Micro-20 development system running Professional OS-9/68020 for initial configuration.

BASIC (Included in PERSONAL PAK)	$ 200.00
C COMPILER (included in PROFESSIONAL PAK)	$ 750.00
PASCAL COMPILER	$ 500.00

UniFLEX (for Micro-20)	$ 400.00
UniFLEX WITH REAL-TIME ENHANCEMENTS	$ 800.00
UniFLEX VM (for TWINGLE-20)	$ 600.00
UniFLEX VM REAL-TIME ENHANCEMENTS	$1000.00

Other Software for UniFLEX

UniFLEX BASIC W/PRECOMPILER	$ 300.00
UniFLEX C COMPILER	$ 350.00
UniFLEX COBOL COMPILER	$ 750.00
UniFLEX SCREEN EDITOR	$ 150.00
UniFLEX TEXT PROCESSOR	$ 200.00
UniFLEX SORT/MERGE PACKAGE	$ 200.00
UniFLEX VSAM MODULE	$ 100.00
UniFLEX UTILITIES PACKAGE I	$ 200.00
UniFLEX PARTIAL SOURCE LICENSE	$1000.00

GMX EXCLUSIVE VERSIONS, CUSTOMIZED FOR THE MICRO-20, OF THE BELOW LANGUAGES AND SOFTWARE ARE ALSO AVAILABLE FROM GMX.

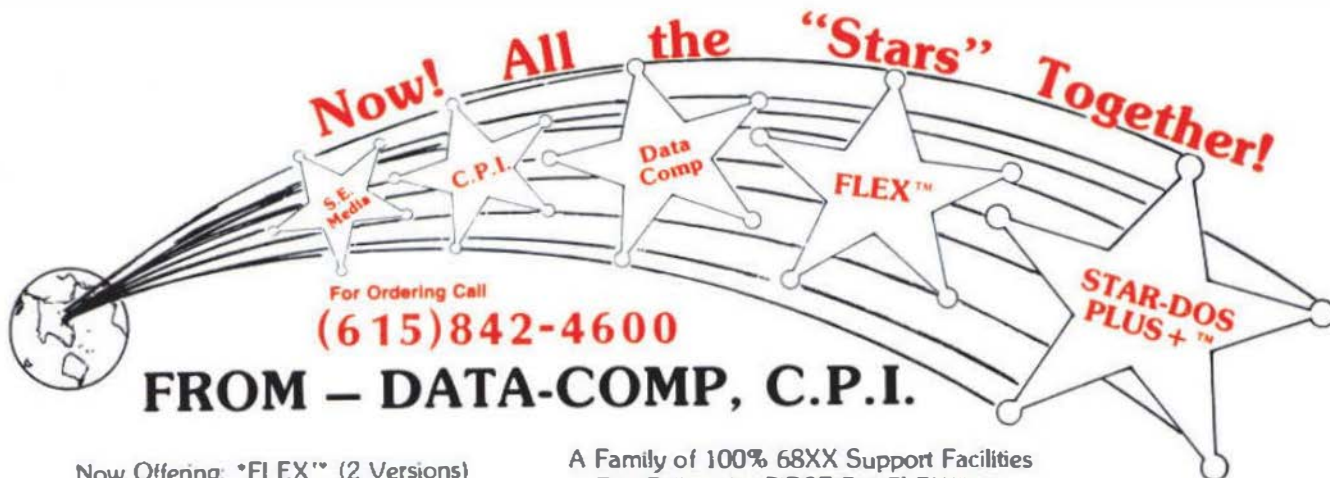|	
ABSOFT FORTRAN (UniFLEX)	$1500.00
SCULPTOR (specify UniFLEX or OS9)	$ 995.00
FORTH (OS9)	$ 595.00
DYNACALC (specify UniFLEX or OS9)	$ 300.00

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ALL PRICES ARE F.O.B. CHICAGO IN U.S. FUNDS

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CONTACT GMX FOR MORE INFORMATION ON THE ABOVE PRODUCTS

GMX STILL SELLS GIMIX S50 BUS SYSTEMS, BOARDS & PARTS. CONTACT GMX FOR COMPLETE PRICE LIST.

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